### $\alpha, \beta$ Operators On Quasi-Boolean Chain

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Abstract:At first,we propose the  $\alpha$  operator and  $\beta$  operator on quasi-Boolean chain.Next investigate their properties and apply them to solve inqualities in one unknown.

keywords:Lattice,Boolean Algebra,Quasi-Boolean Chain,Operator

#### 1 Basic Definitions

Definition 1.1. A partially ordered set (or briefly, a poset) is a pair  $(P, \leq)$  where P is a non-empty set and  $\leq$  is a binary relation on P satisfying for  $x, y, z \in P$ 

- 1.  $x \le x$  (reflexive law)
- 2.  $x \le y, y \le z \Longrightarrow x \le z$  (transitive law)
- 3.  $x \le y, y \le x \Longrightarrow x = y$  (anti-symmetric law)

Where no confusion is likely to arise, it is customary to also use the symbol P to denote  $(P, \leq)$  . Also, if  $x \leq y$  but  $x \neq y$  then we write x < y

Definition 1.2.If P is a poset then a non-empty subset C of P is called a chain in P if and only if any two elements in C are comparable. If P=C then P is a chain. (chains are also referred to as linearly or totally ordered posets) Two members a, b of poset are comparable if  $a \leq b$  or  $b \leq a$ .

Let P be a poset. There is at most one element b in P with the property that  $x \leq b$  for all  $x \in P$ . This element, if it exists, is called the greatest element or the unit

of P. Dually, a least element b(or zero)by the property that  $b \leq x$  for  $x \in P$ . The zero and unit of P are denoted, when they exist, by  $0_p$  and  $1_p$  (or simply by 0,1).

Definition 1.3. A lattice is a poset in which x+y and x.y exist for any  $x,y \in L$ 

In a lattice L, the following two statements are equivalent:

1. 
$$x(y+z) = xy + xz$$
 for all x,y,z in L (1.1)

2. 
$$u + vw = (u + v)(u + w)$$
 for all u,v,w in L (1.2)

Definitin 1.4. A lattice L is distributive if it satisfies one(and hence both)of (1.1) and (1.2)

Definition 1.5. A quasi-Boolean algebra is an algebra of the form (L, (+, ., -)) where + and  $\cdot$  are binary operations, - is a unary operation, satisfying

- 1. (L, (+, .)) is a distributive lattice with 0,1
- 2.  $\overline{x+y} = \overline{x}.\overline{y}; \overline{xy} = \overline{x} + \overline{y}$  for all  $x, y \in L$
- 3.  $\overline{\overline{x}} = x$ . for each  $x \in L$

In this paper, we assume that L is a chain with a kernel element denoted by e, which is defined by the property:  $e=\overline{e}$ 

Example 1:  $L = [0,1], a+b = \min(a,b), \overline{a} = 1-a$ , then (L,(+,.,-)) is a quasi-Boolean algebra, and L is a chain with a kernel element e, where e=0.5

Example 2.  $L = \{a_1, a_2, a_3, a_4, a_5\}, \overline{a_1} = a_5, \overline{a_2} = a_4, a_3 = e, L$  is a quasi-Boolean chain with a kernel element e shown as Fig.1

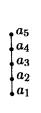


Fig.1

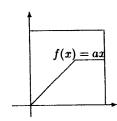
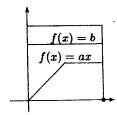


Fig.2



If  $a \leq b$ Then  $S = [0, 1], a\alpha b = 1$ Fig. 3

## 2 Operators on Quasi-Boolean Chain

Let  $a, b \in C$ , define

$$a\alpha b = \sum \{x | ax \le b\}$$

Since C is a complete lattice, so  $\sum \{x | ax \leq b\} \in C$ 

Theorem 2.1.  $a\alpha b$  is the greatest element of the set  $\{x|ax \leq b\}$ 

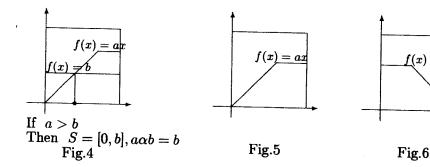
Proof. First,  $a\alpha b \in \{x | ax \leq b\}$ . In fact:

Let  $X(a,b)=\{x|ax\leq b\}$ .  $\forall x_0\in X(a,b),$  then  $ax_0\leq b.$  Now  $\sum\limits_{x_0\in X(a,b)}ax_0=a\sum\limits_{x_0\in X(a,b)}x_0=a(a\alpha b).$  Again,  $\sum\limits_{x_0\in X(a,b)}ax_0\leq b.$  So  $a(a\alpha b)\leq b.$  Hence  $a\alpha b\in X(a,b).$ 

Next. Suppose  $u\in X(a,b)$ , then it is easy to see that  $u\leq \sum_{x_0\in X(a,b)}x_0=a\alpha b$ . So The proof. is completed

Function f(x) = ax as shown in Fig.2

Consider the solution sets of the inequality  $ax \le b$ : (1)  $a \le b$ : The solution set shown as Fig.3;(2) a > b: the solution set shown as Fig.4



Hence we can define a operator  $\alpha$  by the following

$$a\alpha b = \begin{cases} 1, & a \le b \\ b, & a > b \end{cases}$$

Let  $a, b \in C$ . define

$$a\beta b = \sum \{y|ay \le b\overline{x}\}$$

Since C is a complete lattice, so  $~\sum \{y|ay \leq b\overline{y}\} \in C$ 

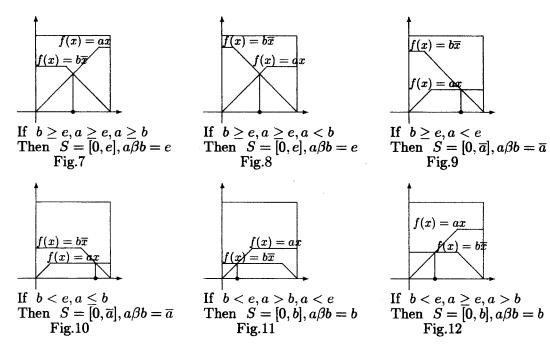
Throrem 2.2.  $a\beta b$  is the greatest element of the set  $\{y|ay \leq b\overline{y}\}$ 

Proof.First,  $a\beta b \in \{y|ay \leq b\overline{y}\}$  Indeed

Let  $Y(a,b)=\{y|ay\leq b\overline{y}\}$ . For each  $p\in Y(a,b)$ , then  $ap\leq b\overline{p}$ ; For any  $q\in Y(a,b)$ , so  $aq\leq b\overline{q}$ . Case 1.  $p\leq q$ : It is obviously that  $ap\leq aq\leq b\overline{q}$ . We can obtain that  $ap\leq \prod\limits_{q\in Y(a,b)}b\overline{q}$ , and than  $ap\leq b\prod\limits_{q\in Y(a,b)}\overline{q}$ . Again  $\prod\limits_{q\in Y(a,b)}\overline{q}=\overline{\sum\limits_{q\in Y(a,b)}q=\overline{a\beta b}}$ . So  $ap\leq b\overline{a\beta b}$ . It is easy to see that  $\sum\limits_{p\in Y(a,b)}ap\leq b\overline{a\beta b}$ . Again  $\sum\limits_{p\in Y(a,b)}ap=a\sum\limits_{p\in Y(a,b)}p=a(a\beta b)$ . Hence  $a(a\beta b)\leq b\overline{a\beta b}$ . Case 2p>q: It is similarly prove that  $a(a\beta b)\leq b\overline{a\beta b}$ .

Next, Suppose  $u \in \{y | ay \le b\overline{y}\}$ . It is easy to see that  $u \le a\beta b$ .

The function  $f_1(x) = ax$ ,  $f_2(x) = b\overline{x}$  shown as Fig.5,Fig.6 respectively. The solution sets of inequality  $ax \le b\overline{x}$  shown as Fig.7—Fig.12,respectively



We can define a operator by the following

$$a\beta b = \begin{cases} e, & b \ge e, a \ge e \\ \overline{a}, & b \ge e, a < e \\ \overline{a}, & b < e, a \le b \\ b, & b < e, a > b \end{cases}$$

It is easy to obtain the following results.

Theorem 2.3.  $\forall a, b \in C$ 

1.  $a\alpha b \geq b$ 

- 2. If  $a \leq b$ , then  $a\beta b \geq e$ .
- 3. If  $e \leq b$ , then  $a\beta b > e$ .

Theovem2.4  $\forall a,b,c \in c$ , then the following statements are equivalent:

- 1.  $ace \leq b$
- 2.  $\overline{a\alpha b} \leq c\alpha b$
- 3.  $\overline{a\alpha b} \leq c\beta b$
- 4.  $\overline{a\beta b} \leq c\alpha b$
- 5.  $\overline{a\beta b} < c\beta b$

Proof. (1)  $\Longrightarrow$  (2) If  $ace \leq b$ , then  $a \leq b$  or  $c \leq b$  or  $e \leq b$ . Case  $a \leq b$ : since  $a\alpha b = 1$ , so  $\overline{a\alpha b} = 0$ . It is obviously that  $\overline{a\alpha b} \leq c\alpha b$ ; Case accupanterise cab = 1, hence  $\overline{a\alpha b} \leq c\alpha b$ ; Case accupanterise caccupanterise caccupanteris

- (2)  $\Longrightarrow$  (1) Suppose that ace > b. We have a > b, c > b, and e > b. By definition, then  $a\alpha b = b$ , and  $c\alpha b = b$ . Again  $\overline{a\alpha b} = \overline{b}$  and e > b, so  $\overline{a\alpha b} > c\alpha b$ . A contradiction.
- $(1) \Longrightarrow (3) \text{ If } ace \leq b, \text{ then } a \leq b \text{ or } c \leq b \text{ or } e \leq b. \text{ case } 1 \text{ } a \leq b \text{ : since } a\alpha b = 1, \text{ so } \overline{a\alpha b} = 0. \text{ It is easy to see that } \overline{a\alpha b} \leq c\beta b; \text{ Case } 2 \text{ } c \leq b \text{ : If } b < e, \text{ then } c\beta b = \overline{c}. \text{ Again } a\alpha b \geq b. \text{ So } \overline{a\alpha b} \leq \overline{b} \leq \overline{c} = c\beta b. \text{ Hence } \overline{a\alpha b} \leq c\beta b; \text{ If } b \geq e, \text{ then } \text{Note that } a\alpha b \geq b, \text{ so } \overline{a\alpha b} \leq \overline{b} < e \leq c\beta b. \text{ Hence } \overline{a\alpha b} \leq c\beta b; \text{ Case } 3 \text{ } e \leq b \text{ : Then } a\alpha b \geq b \text{ and } c\beta b \geq e, \text{ so } \overline{a\alpha b} \leq \overline{b} \leq e \leq c\beta b. \text{ We have } \overline{a\alpha b} \leq c\beta b.$
- (3)  $\Longrightarrow$  (1) Suppose that ace > b, we obtain that a > b, c > b, e > b. So  $a\alpha b = b$  and  $c\beta b = b$ . Then  $\overline{a\alpha b} = \overline{b} > b = c\beta b$ . It follows that  $\overline{a\alpha b} > c\beta b$ . A contradiction.
- $(1) \Longrightarrow (4) \text{ If } ace \leq b, \text{ then } a \leq b \text{ or } c \leq b \text{ or } e \leq b. \text{ Case } 1 \text{ } a \leq b \text{ : If } b < e, \text{ then } a\beta b = \overline{a}. \text{ So } \overline{a\beta b} = a \leq b \leq c\alpha b; \text{ If } b \geq e, \text{ then } a\beta b \geq e. \text{ So } \overline{a\beta b} \leq e \leq b \leq c\alpha b; \text{ Case } 2 \text{ } c \leq b \text{ : we obtain that } c\alpha b = 1. \text{ So } \overline{a\beta b} \leq c\alpha b; \text{ Case } 3 \text{ } e \leq b \text{ : Then } a\beta b \geq e, \text{ so } \overline{a\beta b} \leq e < b \leq c\alpha b. \text{ Hence } \overline{a\beta b} \leq c\alpha b.$
- (4)  $\Longrightarrow$  (1) Suppose that ace > b, then a > b, c > b, e > b. It is easy to see that  $c\alpha b = b, a\beta b = b$ . So  $\overline{a\beta b} = \overline{b} > b = c\alpha b$ . A contradiction.

- $(1) \Longrightarrow (5) \text{ If } ace \leq b, \text{ then } a \leq b \text{ or } c \leq b \text{ or } e \leq b. \text{ Case } 1 \text{ } a \leq b \text{ : If } b < e, \text{ then } a\beta b = \overline{a}. \text{ So } \overline{a\beta b} = a \leq b \leq c\beta b \text{ : If } b \geq e, \text{ then } a\beta b \geq e \text{ and } c\beta b \geq e. \text{ It } follows \text{ that } \overline{a\beta b} \leq e \leq c\beta b; \text{ Case } 2 \text{ } c \leq b \text{ : We have } c\beta b \geq e. \text{ If } b < e, \text{ then } c\beta b = \overline{c}. \text{ It is obviously that } \overline{c\beta b} = c \leq b \leq c\beta b. \text{ Hence } c\beta b = \overline{c} \geq \overline{b} \geq \overline{a\beta b}; \text{ If } b \geq e, \text{ then } a\beta b \geq e \text{ and } c\beta b \geq e. \text{ Clearly } \overline{a\beta b} \leq c\beta b.$
- (5)  $\Longrightarrow$  (1) Suppose that ace > b, Then a > b, c > b, e > b It follows that  $a\beta b$  and  $c\beta b = b$ . It is obviously that  $\overline{a\beta b} = \overline{b} > b = c\beta b$ . A contradiction.

# 3 Inquality in One Unknown on Quasi-Boolean Chain

We have discussed the properties of " $\alpha$ " operator and " $\beta$ " operator on quasi-Boolean chain, Now we apply it to solve Inqualities in one unknown.

Consider the following inquality.

$$a_1x + c_1\overline{x} \le a_2x + b_2x\overline{x} + c_2\overline{x} \tag{3.1}$$

Now construct the inqualities

$$a_1x + c_1\overline{x} \le a_2x \tag{3.11}$$

$$a_1x + c_1\overline{x} \le b_2x\overline{x} \tag{3.12}$$

$$a_1 x + c_1 \overline{x} \le c_2 \overline{x} \tag{3.13}$$

Theorem3.1 The inequality (3.1) is consistent if and only if  $a_1c_1e \leq a_2 + b_2 + c_2$ , and

- 1. If  $a_1c_1e \leq a_2$ , then the set  $[\overline{c_1\beta a_2}, a_1\alpha a_2]$  is the solution of the inequality (3.1)
- 2. If  $a_1c_1e \leq b_2$ , then the set  $[\overline{c_1\beta b_2}, a_1\beta b_2]$  is the solution of the inquality (3.1)
- 3. If  $a_1c_1e \leq c_2$ , then the set  $[\overline{c_1\alpha c_2}, a_1\beta c_2]$  is the solution of inequality (3.1)

Proof. First Inquality (3.11) is consistent if and only if  $[0, a_1 \alpha a_2] \cap [\overline{c_1 \beta a_2}, 1] \neq \emptyset$ . Again  $\overline{c_1 \beta a_2} \leq a_1 \alpha a_2$  if and only if  $a_1 c_1 e \leq b_2$  and inquality (3.12) is consistent if and only if  $a_1 c_1 e \leq b_2$  and inquality (3.13) is consistent if and only if  $a_1 c_1 e \leq c_2$ . Note that inquality (3.1) is consistent if and only if there is at least a inquality among inquality (3.11),(3.12) and (3.13) is consistent and the solution

set ofinquality (3.1) is the union of their solution set. So inequality (3.1) is consistent if and only if  $a_1c_1e \leq a_2 + b_2 + c_2$ .

Next consider the following inquality

$$a_1x + b_1x\overline{x} + c_1\overline{x} \le a_2x + b_2x\overline{x} + c_2\overline{x} \tag{3.2}$$

Construct two inqualities

$$(a_1+b_1)x+c_1\overline{x} \le a_2x+b_2x\overline{x}+c_2\overline{x} \tag{3.21}$$

$$a_1x + (b_1 + c_1)\overline{x} \le a_2x + b_2x\overline{x} + c_2\overline{x} \tag{3.22}$$

It is obviously that

Therem 3.2. Inequality (3.2) is consistent if and only if  $a_1c_1e \le a_2 + b_2 + c_2$  and its solution set is the union of the solution set of inequalities (3.21),(3.22)

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