

On n-reverse of implicators

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Abstract. For given implicator I we can construct the binary operation $I^* : [0;1]^2 \rightarrow [0;1]$ such that $I^*(x,y) = \max\left[0, \min\left(1, n(x) - n(y) + I(n(x), n(y))\right)\right]$, where n is an negator. The binary operation I^* is called the reverse of implicator I . Generally, I^* can't be an implicator. The implicator I is reversible if I^* is an implicator too. The condition of reversibility of implicators are studied.

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Preliminaries

Definition 1. The unary operator $n : [0;1] \rightarrow [0;1]$ is called *negator* iff

$$(n1) \quad a \leq b \Rightarrow n(a) \geq n(b)$$

$$(n2) \quad n(0) = 1, n(1) = 0$$

The negator n is called *involutive* if $n(n(a)) = a$ for any $a \in [0,1]$

Definition 2. The binary operation $I : [0;1]^2 \rightarrow [0;1]$ is called *implicator* if for any $a, b, c \in [0;1]$ hold

$$(i1) \quad b \leq c \Rightarrow I(a, b) \leq I(a, c), \quad I(b, a) \geq I(c, a) \quad (\text{hybrid monotonicity})$$

$$(i2) \quad I(0,0) = I(1,1) = 1, \quad I(1,0) = 0 \quad (\text{boundary conditions})$$

Corollary 1. Using (i1), (i2) we have: $1 = I(0,y) = I(x,1)$ for all $x, y \in [0,1]$

Definition 3. The implicator I is *contrapositive* with respect to negator n if

$$I(x,y) = I((n(y), n(x))) \text{ for all } x, y \in [0,1]$$

Definition 4. The implicator I is *border* if $I(1,y) = y$ for all $y \in [0,1]$.

There exist many ways how to construct implicators. One of them use t-norms and t-conorms.

Definition 5. The binary operation $T : [0;1]^2 \rightarrow [0;1]$ is called *t-norm* if for any $a, b, c \in [0;1]$ hold

$$(t1) \quad T(1, a) = a \quad (\text{boundary condition})$$

$$(t2) \quad a \leq b \Rightarrow T(a, c) \leq T(b, c) \quad (\text{monotonicity})$$

$$(t3) \quad T(a, b) = T(b, a) \quad (\text{commutativity})$$

$$(t4) \quad T(T(a, b), c) = T(a, T(b, c)) \quad (\text{associativity})$$

The binary operation $S : [0;1]^2 \rightarrow [0;1]$ defined by

$$S(x,y) = n(T(n(x), n(y)))$$

where T is a t-norm and n is an involutive negator, is called *t-conorm* dual to the t-norm T with respect to the negator n . Evidently then $S(x,y) = n(T(n(x), n(y)))$. If $n(x) = 1-x$ we shall simply say that S is dual to T .

Example 1. The family of Frank t-norms $\{T_s^F, s \in [0, \infty]\}$ is defined by

$$T_0^F(x, y) = T_M(x, y) = \min(x, y) \quad T_1^F(x, y) = T_P(x, y) = xy$$

$$T_\infty^F(x, y) = T_L(x, y) = \max(0, x + y - 1)$$

$$T_s^F(x, y) = \log_s(1 + (s^x - 1)(s^y - 1) / (s - 1)), s > 0, s \neq 1$$

The family of Frank t-conorms is given by $S_s^F(x, y) = 1 - T_s^F(1 - x, 1 - y)$

Let T be a t-norm S be a dual t-conorm, n be a negator then the binary operation

$$I(x, y) = n(T(x, n(y))) = S(n(x), y)$$

is implicator. If n is an involutive negator then such implicator I is so-called **model implicator** i.e. contrapositive border and fullfils the exchange principle:

$$I(a, I(b, c)) = I(b, I(a, c))$$

Example 2. Let $T(x, y) = T_P(x, y) = xy$ and $n(x) = 1 - x$. We obtain the Reichenbach implicator: $I_P(x, y) = 1 - x(1 - y) = 1 - x - xy$

Example 3. Let $T(x, y) = T_L(x, y) = \max(0, x + y - 1)$, $n(x) = 1 - x$. We obtain the Lukasiewicz implicator $I_L(x, y) = \min(1, 1 - x + y)$

Example 4. $T(x, y) = T_M(x, y) = \min(x, y)$ and $n(x) = 1 - x$. We obtain the Kleene-Dienes implicator $I_M(x, y) = \max(1 - x, y)$

Now we shall show some way how to generate implicators from given ones.

The reverse of implicators

Definition 6. Let I be an implicator, n be a negator. The binary operation $I_n^*: [0, 1]^2 \rightarrow [0, 1]$ defined by

$$I_n^*(x, y) = /n(x) - n(y) + I(n(x), n(y)) /$$

where $/a/ = \min(1, \max(0, a))$, is called the **n-reverse of implicator I**.

Evidently, $I_n^*(0, 0) = I_n^*(1, 1) = I_n^*(0, 1) = 1$, $I_n^*(1, 0) = 0$, $I_n^*(x, 0) = n(x)$, $I_n^*(1, y) = 1 - n(y)$,

but I^* need not be an implicator. If I_n^* is an implicator then we shall say that I is **n-reversible**.

Example 5. Let $I(x, y) = \begin{cases} 1 & \text{if } x=0 \text{ or } y=1 \\ 0 & \text{otherw.} \end{cases}$, n be a continuous negator. Take $y \in (0, 1)$ such

that $n(y) \in (0, 1)$. Then $I_n^*(0, y) = /1 - n(y) + I(1, n(y)) / = 1 - n(y) \neq 1$ which contradicts (i1).

Thus I_n^* isn't implicator.

Theorem 1. If n is continuous negator and I_n^* is an implicator then $I(1, y) \geq y$.

Proof. If I_n^* is an implicator then $1 = I_n^*(0, y) = /1 - n(y) + I(1, n(y)) /$. It follows that

$1 - n(y) + I(1, n(y)) \geq 1$ or $I(1, n(y)) \geq n(y)$. The continuity of n gives the required inequality.

Let n_1 be another negator. Then we define Lukasiewics n_1 -implicator as follows:

$$I_{n_1, L}(x, y) = \min(1, n_1(x) + y)$$

Theorem 2. Let n be a continuous negator. Then $I_{n_1, L}$ is n-reversible iff

$n_1(y) \geq 1 - y$ for all $y \in [0, 1]$.

Proof. $(I_{n_1, L})_n^* = /n(x) - n(y) + \min(1, n_1(n(x)) + n(y)) /$

If $(I_{n_1, L})_n^*$ is an implicator then $1 = (I_{n_1, L})_n^*(x, 1) = /n(x) + n_1(n(x)) /$.

Thus $n(x) + n_1(n(x)) \geq 1$. The continuity of n gives $y + n_1(y) \geq 1$ Conversely, let $y + n_1(y) \geq 1$.

Then $(I_{n_1, L})_n^*(y, y) = / \min(1, n_1(n(y)) + n(y)) / = 1$. It can be easily shown that

$$(I_{n_1, L})_n^* = 1 \quad \text{for } y > x$$

and

$$(I_{n_1, L})_n^*(x, y) = /n(x) - n(y) + 1 / \quad \text{for } y \geq x$$

Such binary operation is the implicator.

Example 6. Let $n(x) = 1 - x$, $n_1(x) = 1 - \sqrt{2x - x^2}$. Then

$$(I_{n_1, L})_n^* = /y - x + \min(1, 2 - y - \sqrt{1 - x^2}) /$$

isn't implicator. E.g. $(I_{n_1, L})_n^*(0.5; 1) = 0.63$

In the following we shall be using the negator $n(x) = 1 - x$ only. For such negator we shall denote $I^* = I_n^*$. It means

$$I^*(x, y) = /y - x + I(1 - x, 1 - y) /$$

If I^* is the negator then we say that I is reversible.

Theorem 3. If I is reversible implicator then I^* is an border implicator.

Proof. $I^*(1, y) = /y - 1 + I(0, 1 - y) / = y$.

Then Theorem 3 allows to generate border implicators from another ones.

Theorem 4. If I is a contrapositive (with respect to $n(x) = 1 - x$) reversible negator The I^* is contrapositive negator.

Proof. $I^*(1 - y, 1 - x) = /y - x + I(y, x) / = /y - x + I(1 - x, 1 - y) / = I^*(x, y)$

Example 7. Let $I(x, y) = I_L(x, y) = \min(1, 1 - x + y)$ (Lukasiewics implicator) Then $I^* = I_L$.

Theorem 5. If $I \geq I_L$ then I is reversible and $I^* = I_L$

Proof. If $x \leq y$ then $1 - x \geq 1 - y$ and

$$I^*(x, y) = /y - x + I(1 - x, 1 - y) / \geq /y - x + I_L(1 - x, 1 - y) / = 1 = I_L(x, y)$$

If $x > y$ then $1 - x < 1 - y$ and

$$I^*(x, y) = /y - x + I(1 - x, 1 - y) / \geq /y - x + 1 / = I_L(x, y)$$

Theorem 6. If I is a reversible implicator then $I(x, y) \geq I_M(x, y) = \text{Max}(1 - x, y)$.

Proof. If I is reversible then $1 = I^*(x, 1) = /1 - x + I(1 - x, 0) /$. Thus

$I(1 - x, 0) \geq x$ or $I(x, 0) \geq 1 - x$. Similarly, $1 = I^*(0, y) = /y + I(1, 1 - y) /$. Therefore $I(1, 1 - y) \geq 1 - y$ or $I(1, y) \geq y$

The next theorem gives some necessary and sufficient condition for the reversibility of implicators (without proof)..

Theorem 7. Let I be an implicator. Denote

$$Q(x, y) = \min(I(x, y), I_L(x, y)) = \min(1, I(x, y), 1 - x + y).$$

Then I is reversible iff $|Q(x, y) - Q(x, z)| \leq |z - y|$ and $|Q(x, y) - Q(z, y)| \leq |x - z|$ and $I(x, y) \geq I_M(x, y) = \text{Max}(1 - x, y)$

Selfreversibility

If $I = I^*$ then the implicator I is called **selfreversible**. Example 6 implies that Lukasiewics implicator I_L is selfreversible. For given t-norm T we can define a binary operation

$T^* : [0, 1]^2 \rightarrow [0, 1]$ such that

$$T^* = /x + y - 1 + T(1 - x, 1 - y)/ = \max(0, x + y - 1 + T(1 - x, 1 - y))$$

If T^* is t-norm too then T is called the t-reversible t-norm. If $T = T^*$ then T is called selfreversible. The selfreversibility of t-norm T allows to generate a selfreversible implicator.

Theorem 8. Let T be a selfreversible t-norm, S is dual t-norm then the implicator

$$I(x, y) = 1 - T(x, 1 - y) = S(1 - x, y)$$

is the selfreversible implicator.

Proof.

Remark. In [] was proved that t-norm T is selfreversible iff T is a Frank norm or a symmetrical ordinal sum of Frank t-norm. Recall that the ordinal sum of t-norm is the t-norm defined as follows:

Let $\{(a_k, b_k), k \in K\}$ be a family of pairwise disjoint subintervals of $[0, 1]$ and $\{T_k, k \in K\}$ be a family of t-norms different from T_M . Then The ordinal sum $\{a_k, b_k, T_k\}$ is the t-norm defined by

$$T(x, y) = a_k + (b_k - a_k)T\left(\frac{x - a_k}{b_k - a_k}, \frac{y - a_k}{b_k - a_k}\right) \text{ if } x, y \in [a_k, b_k]$$

$$T(x, y) = \min(x, y) \text{ otherwise}$$

The ordinal sum $\{a_k, b_k, T_k\}$ is symmetrical if for any $i \in K$ there exists $j \in K$ such that

$$a_i + b_j = 1, a_j + b_i = 1 \text{ and } T_i = T_j$$

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