

ANOTHER PROOF OF FUZZY POSYNOMIAL GEOMETRIC PROGRAMMING DUAL THEOREM

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Abstract Based on the theorem of type fuzzy Kuhn-Tucker of fuzzy convex programming, a dual theorem is obtained for fuzzy posynomial geometric programming. Another proof is given in this paper before the algorithm is obtained for a dual programming.

Keywords Fuzzy posynomial. Geometric programming. Dual theorem. Proof.

Let a prime fuzzy posynomial geometric programming be

$$\begin{aligned} \min \quad & \underline{g}_0(x), \\ \text{s.t.} \quad & \underline{g}_m(x) \leq 1, (m=1, 2, \dots, M) \\ & x > 0 \end{aligned} \quad (1)$$

and a nonnegative T-vector be given tallying with $\sum_{i=1}^{T_n} a_{m,i} = 1$ ($m=0, 1, \dots, M$), $a = (\delta_{01}, \dots, \delta_{0T_0}, \dots, \delta_{M1}, \dots, \delta_{MT_M})^T \geq 0$. Let

$$\bar{g}_m(x) = \sum_{i=1}^{T_n} \left(\frac{v_{m,i}(x)}{\delta_{m,i}} \right)^{a_{m,i}} = \sum_{i=1}^{T_n} \left(\frac{c_{m,i}}{\delta_{m,i}} \prod_{n=1}^N x_n^{\gamma_{m,i,n}} \right)^{a_{m,i}} = \bar{c}_m \prod_{n=1}^N x_n^{\bar{\gamma}_n} \quad (2)$$

where $\bar{c}_m = \prod_{i=1}^{T_n} \left(\frac{c_{m,i}}{\delta_{m,i}} \right)^{a_{m,i}}$, $\bar{\gamma}_n = \sum_{i=1}^{T_n} \gamma_{m,i,n} \delta_{m,i}$.

If $\bar{g}_m(x)$ is used instead of $\underline{g}_m(x)$ ($m=0, 1, \dots, M$), then a single fuzzy geometric programming is obtained as follows.

$$\begin{aligned} (P_0) \quad & \min \quad \bar{g}_0(x) \\ \text{s.t.} \quad & \bar{g}_m(x) \leq 1, (m=1, 2, \dots, M) \\ & x > 0 \end{aligned}$$

while the dual form of (1) is

$$\begin{aligned} \max \quad & \underline{d}(\omega) \\ \text{s.t.} \quad & \underline{\Gamma}^T \omega = 0, \alpha, \beta \in [0, 1] \\ & \omega \geq 0 \end{aligned} \quad (3)$$

where $\underline{d}(\omega) = (a_0/\omega_{00})^{\alpha_0} \prod_{i=1}^{T_0} (c_{0i} a_{0i}/\omega_{0i})^{\alpha_i} \prod_{m=1}^M \prod_{i=1}^{T_m} (c_{mi} a_{mi}/\omega_{mi})^{\alpha_i} \omega_{m0}^{\alpha_m}$.

Definition 1 M_P and M_D denote the constraint least upper and greatest lower bounds of (1) and (3) respectively, i.e.

$$\begin{aligned} M_P &= \inf \underline{g}_0(x) & M_D &= \sup \underline{d}(\omega) \\ \text{s.t.} \quad & \underline{g}_m(x) \leq 1 (m=1, \dots, M), & \text{s.t.} \quad & \omega_{00} = 1, \sum_{i=1}^{T_n} \gamma_{0i,n} \omega_{0i} = \omega_{00} (n=1, \dots, N), \end{aligned}$$

Lemma 1 If for continuous nondecreasing functions (CMNF) \underline{B} , \underline{c} and $\underline{\gamma}$ with $\underline{g}_0(x^*) = \underline{d}(x^*)$, then x^* and ω^* must be an optimal solution of (1) and (3) respectively, and for any fuzzy feasible solution x in [3], we have

$$\underline{g}_0(x^*) \prod_{m=0}^M \prod_{t=1}^{T_m} (\omega_{m,t}^*)^{\alpha_{m,t}} = \prod_{m=0}^M (\omega_{m,0}^*)^{\alpha_{m,0}} \prod_{m=0}^M \underline{c}_{m,t}^{\alpha_{m,t}} \quad (4)$$

Proof For any fuzzy feasible solution x of (1), from Lemma 4.1 in [2], we have

$$\underline{g}_0(x) \geq \underline{d}(\omega^*) = \underline{g}_0(x^*).$$

Therefore x^* must be a fuzzy optimal solution of (1). We can prove in the similar way that x^* and ω^* must tally with

$$\omega_{m,t} = \begin{cases} \underline{v}_0(x)/\underline{g}_0(x), & (m=0, t=1, \dots, T_0), \\ \omega_{m,0} \underline{v}_{m,t}(x), & (m \neq 0; t=1, \dots, T_m). \end{cases}$$

Therefore $(\underline{g}_0(x^*) \omega_{0,t}^*) = (\underline{v}_0(x^*))^{\alpha_{0,t}}$, $(\omega_{m,t}^*)^{\alpha_{m,t}} = (\omega_{m,0}^* \underline{v}_{m,t}(x^*))^{\alpha_{m,t}}$ ($m \neq 0$), from $\omega_{0,0} = 1$, we have

$$\begin{aligned} \underline{g}_0(x^*) \prod_{m=0}^M \prod_{t=1}^{T_m} (\omega_{m,t}^*)^{\alpha_{m,t}} &= \prod_{m=0}^M (\omega_{m,0}^*)^{\alpha_{m,0}} \prod_{m=0}^M \prod_{t=1}^{T_m} (\underline{v}_{m,t}(x^*))^{\alpha_{m,t}} \\ &= \prod_{m=0}^M (\omega_{m,0}^*)^{\alpha_{m,0}} \prod_{m=0}^M \prod_{t=1}^{T_m} \underline{c}_{m,t}^{\alpha_{m,t}} \prod_{n=1}^N x_n^{\sum_{m=0}^M \sum_{t=1}^{T_m} \gamma_{m,t,n} \alpha_{m,t}} (\Gamma^T \omega = \sum_{m=0}^M \sum_{t=1}^{T_m} \gamma_{m,t,n} \omega_{m,t}), \end{aligned}$$

and from orthogonality condition, Eq. (4) is known to be true.

Lemma 2 For a CMNF \underline{B} , \underline{c} and $\underline{\gamma}$, when (1) and (3) are both fuzzy consistent, there must be

$$0 < M_D \leq M_P < \infty.$$

Proof Since (3) is fuzzy consistent, $M_D > 0$, and (1) fuzzy consistent, $M_P < \infty$

$$\underline{g}_0(x) \geq \underline{g}_0(x) \prod_{m=0}^M (\underline{g}_m(x))^{\alpha_m} \geq \underline{d}(\omega). \quad [1] [2]$$

Therefore, the lemma holds.

Theorem 1 Let \underline{B} , \underline{c} and $\underline{\gamma}$ be a CMNF. If x is a fuzzy feasible solution of (1), it must be that of a single posynomial fuzzy geometric programming (\underline{P}_1). If (1) is fuzzy consistent, so is (\underline{P}_1). Again $M_D \leq M_P$, here M_D and M_P denote the constrained greatest lower bound of (1) and (\underline{P}_1) respectively. Then, the dual form of (\underline{P}_1) is

$$\begin{aligned} (D_1) \quad & \max (\underline{a}_0' / \omega_0')^{\alpha_0'} \prod_{m=0}^M (\underline{c}_m \underline{a}_m / \omega_m)^{\alpha_m} \\ \text{s.t.} \quad & \omega_0' = 1, \bar{\gamma}_0^* \omega_0 = \omega_0', \quad \bar{\gamma}_0^* \text{ is an index of } \underline{g}_0(x). \\ & \sum_{m=1}^M \bar{\gamma}_m \omega_m = 0, \quad (n=1, \dots, N), \quad \alpha, \beta \in [0, 1], \\ & \omega \geq 0. \end{aligned}$$

Proof If we use $\bar{\underline{g}}_m(x)$ instead of $\underline{g}_m(x)$ ($m=0, 1, \dots, M$), then (1) is turned into (\underline{P}_1). Owing to $(\sum_{i=1}^T \underline{v}_i)^{\alpha_i} \geq \prod_{i=1}^T (\underline{v}_i / \omega_i)^{\alpha_i} \omega_0^{\alpha_0}$ with $\omega_i / \omega_0 = \varepsilon_{m,i}$, we use $\varepsilon_{m,i}$ as weight vector, from

$$\implies \prod_{t=1}^{T_n} (v_{m_t}(x)/\varepsilon_{m_t}\omega_0)^{\varepsilon_{m_t}} = \prod_{t=1}^{T_n} (v_{m_t}(x)/\varepsilon_{m_t})^{\varepsilon_{m_t}} = \bar{c}_m \prod_{n=1}^N x_n^{\bar{z}_n} \leq \sum_{t=1}^{T_n} v_{m_t}(x).$$

Obviously, $\bar{g}_m(x) \leq g_m(x)$, such that a fuzzy feasible solution is that of (P_1) .

$$(P_1) \iff \begin{aligned} & \min \bar{c}_0 \prod_{n=1}^N x_n^{\bar{z}_n} && \max ((1-\alpha)/\bar{\omega}'_0)^{\bar{\alpha}} \prod_{m=0}^M (c_m^{-1}(\beta) B_m^{-1}(\alpha)/\bar{\omega}_m)^{\bar{\omega}_m} \\ & \text{s.t. } \bar{c}_m \prod_{n=1}^N x_n^{\bar{z}_n} \leq 1 && \text{s.t. } \bar{\omega}'_0 = 1, \bar{\gamma}_0^{-1}(\beta) \bar{\omega}_0 = \bar{\omega}'_0, \alpha, \beta \in [0, 1] \\ & x > 0 && \sum_{m=0}^M \bar{\gamma}_m^{-1}(\beta) \bar{\omega}_m = 0, \bar{\omega} \geq 0 (n=1, \dots, N) \end{aligned} \iff (\bar{D}_1)$$

where $\bar{\omega} = (\bar{\omega}_0, \bar{\omega}_1, \dots, \bar{\omega}_M)^T$ is $(m+1)$ -dimension dual parameter vector.

Theorem 2 Let $B_m, B_m^{-1}; c_m, c_m^{-1}$ and γ_m, γ_m^{-1} be continuous and strictly monotonous. If (P_1) is fuzzy consistent with $M_{P_1} > 0$, then there must exist dual feasible solution ω' , such that weight vector δ ,

$$\varepsilon_{m_t} \omega'_{m_0} = \omega'_{m_t} \quad (m=0, 1, \dots, M; t=1, 2, \dots, T_m) \tag{5}$$

with $M_{P_1} = d(\omega')$.

Conversely, if dual fuzzy feasible solution ω' makes (5) hold, $M_{P_1} > 0$ must exist.

Proof Owing to $(P_1) \iff$

$$\begin{aligned} & \max(1-\alpha) && \max(1-\alpha) \\ & \text{s.t. } B_m(c_m^{-1}(\beta)) \prod_{n=1}^N x_n^{\bar{z}_n^{-1}(\alpha)} \geq \alpha, x > 0, && \text{s.t. } \bar{c}_m^{-1}(\beta) B_m^{-1}(\alpha) \prod_{n=1}^N x_n^{\bar{z}_n^{-1}(\alpha)} \leq 1, x > 0, \quad (\bar{P}_1) \\ & \alpha, \beta \in [0, 1], m=0, 1, \dots, M && \alpha, \beta \in [0, 1], m=0, 1, \dots, M \\ & \iff \begin{aligned} & \max(1-\alpha) \\ & \text{s.t. } \sum_{n=1}^N \bar{\gamma}_m^{-1}(\beta) z_n \leq \log(\bar{c}_m^{-1} B_m^{-1}(\alpha)) \\ & (m=0, 1, \dots, M), \alpha, \beta \in [0, 1]. \end{aligned} \end{aligned}$$

If Let $M_{\bar{P}_1} > 0$ as the proof in Theorem 1.7.2 in [3], then its linear programming (\bar{P}_1) has the greatest lower bound $\log M_{\bar{P}_1}$, hence, from dual theory of an ordinary linear programming, a linear one corresponding to (\bar{D}_1) :

$$\begin{aligned} & \max \{ \bar{\omega}'_0 \log(1-\alpha) + \sum_{m=0}^M \bar{\omega}_m \log \bar{c}_m^{-1}(\beta) B_m^{-1}(\alpha) \} \\ & \text{s.t. } \bar{\omega}'_0 = 1, \bar{\gamma}_0^{-1}(\beta) \bar{\omega}'_0 = 1 \\ & \sum_{m=0}^M \bar{\gamma}_m^{-1}(\beta) \bar{\omega}_m = 0, \alpha, \beta \in [0, 1], (n=1, 2, \dots, N) \end{aligned} \tag{6}$$

has the greatest value $\log M_{P_1}$. Moreover (\bar{D}_1) has fuzzy biggest value M_{P_1} . Now let $\bar{\omega}^{*'}$ be an optimal solution of (\bar{D}_1) and (6), then

$$M_{P_1} = ((1-\alpha)/\bar{\varepsilon}_{0,t})^{\bar{\omega}'_0} \prod_{m=0}^M \prod_{t=1}^{T_m} (c_{m_t}^{-1}(\beta) B_{m_t}^{-1}(\alpha)/\bar{\varepsilon}_{m_t})^{\bar{\omega}_m}$$

and from

$$\bar{\varepsilon}_{0,t} = \bar{\omega}'_{0,t} / \bar{\omega}'_0, \quad \bar{\varepsilon}_{m_t} = \bar{\omega}'_{m_t} / \bar{\omega}'_m, \quad \bar{\omega}'_{m_0} = \sum_{t=1}^{T_m} \bar{\omega}'_{m_t} = \sum_{t=1}^{T_m} \bar{\varepsilon}_{m_t} \bar{\omega}'_m = \bar{\omega}'_m \tag{7}$$

Therefore $M_{\bar{P}_i} = ((1-\alpha)\bar{\omega}_{00}^* / \bar{\omega}'_0)^{\bar{\omega}_0} \prod_{m=0}^M \prod_{t=1}^{T_m} (c_{m,t}^{-1}(\beta) B_m^{-1}(\alpha) \bar{\omega}_{m0}^* / \bar{\omega}_{m,t}^*)^{\bar{\omega}_m} = d(\bar{\omega}^*)$ (8)

Since $\bar{\omega}^*$ is a dual feasible solution of (\bar{D}_i) , and from Eq. (2) and (7), we have

$$\sum_{m=0}^M \sum_{t=1}^{T_m} \gamma_{m,t}^{-1}(\beta) \bar{\omega}_{m,t}^* = \sum_{m=0}^M \sum_{t=1}^{T_m} \gamma_{m,t}^{-1}(\beta) \bar{\varepsilon}_{m,t} \bar{\omega}'_m = \sum_{m=0}^M \gamma_{m,t}^{-1}(\beta) \bar{\omega}'_m = 0.$$

Again since $\bar{\omega}' \geq 0$, $\bar{\varepsilon} \geq 0$, then $\bar{\omega}^* \geq 0$ with $\bar{\omega}_{00}^* = \bar{\omega}'_0 = 1$. Hence $\bar{\omega}^*$ is certainly a feasible solution of (\bar{D}_i) , moreover ω^* is a fuzzy feasible one of (3).

Conversely, let $\bar{\omega}^*$ be a dual feasible solution which makes (7) hold, and define $\bar{\omega}'_m = \bar{\omega}_{m0}^*$ ($m=0, 1, \dots, M$), we have $\bar{\omega}'_0 = 1$ and $\bar{\omega}'_m \geq 0$ with

$$\sum_{m=0}^M \gamma_{m,t}^{-1}(\beta) \bar{\omega}'_m = \sum_{m=0}^M \sum_{t=1}^{T_m} \gamma_{m,t}^{-1}(\beta) \bar{\varepsilon}_{m,t} \bar{\omega}'_m = \sum_{m=0}^M \sum_{t=1}^{T_m} \gamma_{m,t}^{-1}(\beta) \bar{\omega}_{m,t}^*.$$

Therefore, $\bar{\omega}'$ is a feasible solution of (D_i) from Theorem 1.7.2 in [3], and $M_{P_i} > 0$.

Since (P_i) consistence $\iff (\bar{P}_i)$ consistence $\iff (\bar{D}_i)$ consistence. While

$$(7) \iff \underline{B}(\bar{\varepsilon}_{m,t}, \bar{\omega}_{m0}^*) = \underline{B}(\bar{\omega}_{m,t}^*), \text{ i.e. } \varepsilon_{m,t} \omega_{m0}^* = \omega_{m,t}^*,$$

$$(8) \iff \underline{B}(M_{\bar{P}_i}) = \underline{B}(d(\bar{\omega}_{m,t}^*)), \text{ i.e. } M_{\bar{P}_i} = d(\omega^*),$$

$$(m=0, 1, \dots, M; t=1, 2, \dots, T_m).$$

$$M_{\bar{P}_i} > 0 \iff M_{P_i} > 0.$$

Therefore, the theorem holds.

Theorem 3 Suppose (1) is fuzzy super consistent with a minimal point x^* , when \underline{B} , \underline{c} and $\underline{\gamma}$ are continuous and strictly monotonous, x^* is a minimal one of single posynomial fuzzy geometric programming (P_i) corresponding to the weight

$$\varepsilon_{m,t} = v_{m,t}(x^*) / g_{m,t}(x^*) (m=0, 1, \dots, M; t=1, 2, \dots, T_m) \quad (9)$$

and $M_{P_i} = M_{P_i}$.

Proof Because of that and according to

$$(1) \iff \begin{aligned} & \max(1-\alpha) \\ & \text{s.t. } \sum_{t=1}^{T_m} c_{m,t}^{-1}(\beta) B_m^{-1}(\alpha) \prod_{n=1}^N x_n^{\gamma_{m,t}(\beta)} \leq 1 \\ & x > 0, \alpha, \beta \in [0, 1] (m=0, 1, \dots, M) \end{aligned} \quad (10)$$

we know super consistence (10) has a minimal point \bar{x}^* . And from the knowledge of Theorem 1.7.3 in [3], a minimal one \bar{x}^* exists in single posynomial geometric programming (\bar{P}_i) corresponding to weight $\varepsilon_{m,t} = v_{m,t}(\bar{x}^*) / g_{m,t}(\bar{x}^*) (m=0, 1, \dots, M; t=1, 2, \dots, T_m)$, and $M_{\bar{P}_i} = M_{\bar{P}_i}$, such that

$$\varepsilon_{m,t} = v_{m,t}(\bar{x}^*) / g_{m,t}(\bar{x}^*), M_{\bar{P}_i} = M_{\bar{P}_i}.$$

Hence the theorem is true.

Theorem 4 Suppose (1) is a fuzzy super consistence with a fuzzy optimal solution x^* , then there must exist Lagrange multiplier $\mu^* = (\mu_1^*, \mu_2^*, \dots, \mu_h^*)^T \geq 0$ with continuous and strictly nondecreasing functions \underline{B} , \underline{c} and $\underline{\gamma}$, such that

$$\nabla g_0(x^*) + \sum_{m=1}^M \mu_m^* \nabla g_m(x^*) = 0$$

while ω^* defined by

$$\omega_{m,t}^* = \begin{cases} v_{0,t}(x^*)/\underline{g}_0(x^*) & (m=1, 2, \dots, M) \\ \mu_m^* v_{m,t}(x^*)/\underline{g}_0(x^*) & (m \neq 0, t=1, 2, \dots, T_m) \end{cases}$$

is a dual fuzzy optimal solution of (3). And

$$\underline{g}_0(x^*) = d(\omega^*).$$

Proof First, from Eq. (9), and the knowledge of Theorem 3, a minimal point corresponding to (P_1) is x^* with a minimal value being $\bar{g}_0(x^*) = \underline{g}_0(x^*) > 0$; but from Theorem 2, $\exists: \omega^*$ tallying with (5), such that

$$d(\omega^*) = \bar{g}_0(x^*) = \underline{g}_0(x^*).$$

From Theorem 1, we know, ω^* is a fuzzy optimal solution of (3).

Lastly, from Eq. (5) and (9), then

$$\omega_{m,t}^* = \varepsilon_{m0}^* = v_{m,t}(x^*)\omega_{m0}^* / \underline{g}_m(x^*),$$

when $m=0$, $\omega_{0,t}^* = v_{0,t}(x^*)/\underline{g}_0(x^*)$ owing to $\omega_{00}^* = 1$; when $m \neq 0$, owing to

$$\omega_{m0}^* = \sum_{t=1}^{T_m} \omega_{m,t}^* = \mu_m^* \sum_{t=1}^{T_m} v_{m,t}(x^*)/\underline{g}_0(x^*) = \mu_m^* \underline{g}_m(x^*)/\underline{g}_0(x^*) \stackrel{\text{from(11)}}{=} \mu_m^* / \underline{g}_0(x^*).$$

Hence

$$\omega_{m,t}^* = v_{m,t}(x^*)\mu_m^* / \underline{g}_m(x^*)\underline{g}_0(x^*) = \mu_m^* v_{m,t}(x^*)/\underline{g}_0(x^*).$$

Suming up, we have

$$\omega_{m,t}^* = \begin{cases} v_{0,t}(x^*)/\underline{g}_0(x^*), & (m=0; t=1, \dots, T_0), \\ \mu_m^* v_{m,t}(x^*)/\underline{g}_0(x^*), & (m \neq 0; t=1, \dots, T_m). \end{cases}$$

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