

A NOTE ON (FUZZY) CONFLICTS (AND CONTRADICTION SETS).

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Basic notions of the conflict theory, introduced by Pawlak [1,2], are shortly presented and followed by a fuzzy approach. A post-scriptum on Shi's [3] contradiction sets is given.

Keywords: conflicts theory, fuzzy conflicts, contradiction sets.

Notations: # A cardinality of A set,
 Δ_A diagonal (all (x,x)) of A set.

1. CONFLICTS [2].

Let U be a some finite set with $\#U = n \geq 2$ and $L = \{-1, 0, +1\} \subseteq \mathbb{R}$.

Basic definitions.

By giving a relationship function $\varphi: U \times U \rightarrow L$ a couple

$(x, y) \in U \times U$ is called:

- ally iff $\varphi(x, y) = +1$,
- neutral iff $\varphi(x, y) = 0$,
- enemy iff $\varphi(x, y) = -1$.

The couple $C = (U, \varphi)$ is called the configuration (of U).

By C are definible 3 (disjoint binary) relations on U:

$$R_C^+ = \varphi^{-1}(\{+1\}), R_C^- = \varphi^{-1}(\{-1\}), R_C^0 = \varphi^{-1}(\{0\})$$

called, resp., alliance, conflict and neutrality.

If $R_C^- = \emptyset$ C is conflictless else conflict configuration;

$R_C^+ = U^2$ is called total alliance, $R_C^- = U^2 - \Delta$ total conflict;

$A_C = R_C^+ \cup R_C^-$ is called activity.

Regularity.

The configuration $C = (U, \varphi)$ is called regular iff

$$R1 - \Delta \subseteq R_C^+$$

$$R2 - (R_C^+)^{-1} = R_C^+$$

$$R3 - (R_C^-)^{-1} = R_C^-$$

$$R4 - R_C^+ \circ R_C^+ \subseteq R_C^+$$

$$R5 - R_C^- \circ R_C^- \subseteq R_C^+$$

$$R6 - R_C^- \circ R_C^+ \subseteq R_C^-$$

$$R7 - R_C^+ \circ R_C^+ \subseteq R_C^-$$

(if $n \geq 3$, $R1, R6$ are independents), else non-regular.

Proposition [4].

$C = (U, \varphi)$ is regular iff

- A_C is equivalence relation,

$$- \forall (x, y), (y, z) \in A_C \Rightarrow \varphi(x, y) \cdot \varphi(y, z) = \varphi(x, z).$$

2. EXTENSIONS.

Żakowski [4] extend the class of functions of relationship,

by using as range the real interval $I = [-1, +1] = \{x \in \mathbb{R} \mid -1 \leq x \leq 1\}$.

Let $\phi: U^2 \rightarrow I$ be.

The corresponding same, of 1 sect., definitions are analogously

extended to configuration $C_\phi = (U, \phi)$.

The couple $(x, y) \in U^2$ is called: ally iff $\phi(x, y) > 0$,

neutral iff $\phi(x, y) = 0$ and enemy iff $\phi(x, y) < 0$.

The value $\phi(x, y)$ is called the alliance degree.

The configuration $C_\phi = (U, \phi)$ is called extensively regular iff

1. $\phi(x,x) > 0$
2. $\phi(x,y) > 0 \Rightarrow \phi(y,x) > 0$
3. $\phi(x,y) < 0 \Rightarrow \phi(y,x) < 0$
4. $\phi(x,y) > 0, \phi(y,z) > 0 \Rightarrow \phi(x,z) > 0$
5. $\phi(x,y) < 0, \phi(y,z) < 0 \Rightarrow \phi(x,z) > 0$
6. $\phi(x,y) < 0, \phi(y,z) > 0 \Rightarrow \phi(x,z) < 0$
7. $\phi(x,y) > 0, \phi(y,z) < 0 \Rightarrow \phi(x,z) < 0$.

Let $\sigma : I \rightarrow L$ be such that $\sigma(0) = 0$, $\sigma(x) = \frac{x}{|x|}$ if $x \neq 0$

and $\psi = \phi \circ \sigma$:

$$\begin{array}{ccc} U & \xrightarrow{\psi} & \\ \phi \downarrow & & \\ [-1, 1] & \xrightarrow{\sigma} & \{-1, 0, +1\} \end{array} .$$

Corollary.

C_ϕ is ext.regular iff C is regular.

Corollary.

$\phi(x,y) = 0, \phi(y,z) \neq 0 \Rightarrow \phi(x,z) = 0 \quad \forall x,y,z \in U$.

3. FUZZY APPROACH.

Let U be a nonempty universal set, $\Lambda = [0, 1] = \{x \in \mathbb{R} \mid 0 \leq x \leq 1\}$,

$\theta : U^2 \rightarrow \Lambda$ define a fuzzy relation on U . θ is called:

F1- reflexive if $\theta(x,x) = 1 \quad \forall x \in U$

F2- ε -reflexive if $\theta(x,x) \geq \varepsilon$ (with $\varepsilon \in]0, 1[$)

F3- weakly-reflexive if $\theta(x,y) \leq \theta(x,x) \wedge \theta(y,y) \quad \forall (x,y) \in U^2$

F4- symmetric if $\theta(x,y) = \theta(y,x) \quad \forall (x,y) \in U^2$

F5- max-* transitive if $\theta(x,y) * \theta(y,z) \leq \theta(x,z) \quad \forall (x,y,z) \in U^3$

F6- proximity if θ is F1, F4

F7- similarity if θ is F5, F6.

Let $\mathcal{J} \in]0, 1[= \{r \in \mathbb{R} \mid 0 < r < 1\}$ be a threshold.

A couple $(x, y) \in U^2$ is called: \mathcal{J} -fuzzly ally iff $\theta(x, y) > \mathcal{J}$,

\mathcal{J} -fuzzly neutral iff $\theta(x, y) = \mathcal{J}$ and \mathcal{J} -fuzzly enemy iff

$\theta(x, y) < \mathcal{J}$; $\theta(x, y)$ is called the fuzzy alliance degree between x and y . Analogously to 2 sect. by changing \mathcal{J} to 0, the configuration $C_\theta = (U, \theta)$ is defined f.regular.

Let $\lambda_{\mathcal{J}} : \Lambda \rightarrow L$ be s.t. $\lambda_{\mathcal{J}}(r) = +1$ if $r > \mathcal{J}$, $= 0$ if $r = \mathcal{J}$, $= -1$ if $r < \mathcal{J}$ and let $\varphi = \theta \circ \lambda_{\mathcal{J}}$ be.

Corollary.

(U, θ) is f.regular iff (U, φ) is regular.

Example.

$U = \{x, y, z, t\}$, $\theta \in U \times U$, $\theta(a, b) \in \{1/4, 1/2, 2/3, 3/4, 1\}$:

θ	x	y	z	t
x	1	1/2	1/4	2/3
y	1/2	1	3/4	2/3
z	1/4	3/4	1	2/3
t	2/3	2/3	2/3	1

For $\mathcal{J} = 2/3$ it is:

$$\theta_{\mathcal{J}}^+ = \{(x, x), (y, y), (y, z), (z, y), (z, z), (t, t)\}$$

$$\theta_{\mathcal{J}}^- = \{(x, y), (x, z), (y, x), (z, x)\}$$

$$\theta_{\mathcal{J}}^0 = \{(x, t), (y, t), (z, t), (t, x), (t, y), (t, z)\}$$

$A_{\mathcal{J}} = \theta_{\mathcal{J}}^+ \cup \theta_{\mathcal{J}}^-$ ($U/A_{\mathcal{J}} = \{\{x, y, z\}, \{t\}\}$) and (U, θ) is f.regular.

Instead, for $\mathcal{J} = 1/2$ ($A_{\mathcal{J}}(x, t) = 2/3 > \mathcal{J}$, $A_{\mathcal{J}}(t, z) = 2/3 > \mathcal{J}$, $A_{\mathcal{J}}(x, z) = 1/4 < \mathcal{J}$)

or for $\mathcal{J} = 3/4$ ($A_{\mathcal{J}}(z, x) = 1/4 < \mathcal{J}$, $A_{\mathcal{J}}(x, y) = 1/2 < \mathcal{J}$, $A_{\mathcal{J}}(z, y) = 3/4 = \mathcal{J}$)

(U, θ) is not f.regular.

Moreover; for $\mathcal{J} = 2/3$ θ is not max-min transitivity

($\theta(x, y) \wedge \theta(y, z) = 1/2 \wedge 3/4 = 1/2 > 1/4 = \theta(x, z)$) and for bounded difference \ominus ($a \ominus b = 0 \vee (a + b - 1)$) θ is not max- \ominus transitivity

($\theta(x, t) \ominus \theta(t, z) = 2/3 \ominus 2/3 = 1/3 > 1/4 = \theta(x, z)$).

A fuzzy relation $\pi : U^2 \rightarrow \Lambda$ s.t. $\pi \supseteq \theta$ is called pacification of θ (i.e. $\pi(x,y) \geq \theta(x,y) \quad \forall x,y \in U$).

The max-min transitive closure τ of θ , $\tau(x,y) = \bigvee_{u \in U} \{\theta(x,u) \wedge \theta(u,y)\}$, is a fuzzy pacification of θ .

4. POST-SCRIPTUM ON SHI'S CONTRADICTION SETS [3].

A philosophical approach to the sets is presented by Shi [3,5].

The classic logic and set theory use the dichotomy: true/false, one/zero, yes/no. Shi considers the definition/existence (intension/extension) of a set by a propriety. For Shi a concept (propriety) is (philosophically) proponible only if we have in mind/experience the its interlinked contrasts (ego/non-ego). Essentially, this is a "gnoseology" philosophical question.

For Shi, an unclear concept is not a concept (in essence). He distinguishes [5] the imagination for the vague intension and the concept in essence for the clear intension: the first is the intuition-thought and the second the logic-thought.

To make clear the definition of fuzzyness, he try make clear the definition of contradiction. Against to the contradiction as "contrary quality" (opposition) (life/death, correct/wrong, yin/yang), Shi considers the contradiction as "difference quality" (fast/slow, fat/thin).

For its (philosophical) investigations, he propose a "contradiction set" model. Here, I try to shortly give a (my) mathematical translation of Shi's (mathematically presented) ideas.

Definitions.

A concept, called elaborate topic P, determine a contradiction set $\overset{\ominus}{N}_t$, represented by a t-uple $\{a_d \mid 1 \leq d \leq t \ (t \in \mathbb{N})\}$, called t-kind, and any a_d , called vision, is a t_d -uple $\{a_{dj} \mid 1 \leq j \leq t_d \ (j \in \mathbb{N})\}$.

The t-kind is called:

- standard type if $t = 3$
- two-leaf type if $t = 2$
- one-leaf type if $t = 1$
- extensive type if $t \geq 4$
- empty if $t = 0$.

(For Shi, the Cantor's sets are 1-kind and the Zadeh's sets are extensive type).

Moreover, the t-kind of the contradiction set $\overset{\ominus}{N}_t$ is determined by a mental perspective (vision-angle: method of making contradiction).

By writing $\bar{a}_d = \alpha(a_d)$, let $\alpha: \overset{\ominus}{N}_t \rightarrow [0, 1]$ be s.t. $\alpha(\overset{\ominus}{N}_t) = \{\bar{a}_d\}$ and $\sum_1^t d \bar{a}_d = 1$: the function α is called vision-angle.

Example.

Let $\overset{\ominus}{N}_t = \left\{ \{a_{ij} \mid 1 \leq j \leq t_i\} \mid 1 \leq i \leq t \right\}$ be and $n = \sum_1^t t_i$. Let $\alpha: \overset{\ominus}{N}_t \rightarrow [0, 1]$ be s.t. $\bar{a}_i = \alpha(a_i) = t_i/n$. Then $\sum_1^t \bar{a}_i = 1$: α is a vision angle for $\overset{\ominus}{N}_t$.

Shi's operators.

$$\begin{aligned} \overset{\ominus}{N}_t &= \{a_i\}_t & \overset{\ominus}{M}_t &= \{b_i\}_t : \\ \overset{\ominus}{N}_t \subseteq \overset{\ominus}{M}_t & \text{ iff } a_i \subseteq b_i \quad \forall i \in [1, 2, \dots, t] \subseteq \mathbb{N} \\ \overset{\ominus}{N}_t \cup \overset{\ominus}{M}_t &= \{(a_i \cup b_i)_i\}_t \\ \overset{\ominus}{N}_t \cap \overset{\ominus}{M}_t &= \{(a_i \cap b_i)_i\}_t . \end{aligned}$$

Re-definitions.

A function $a: \Omega \rightarrow \Lambda$ is called an attribute and $a(\omega)$ ($\omega \in \Omega$) the attribute value in ω .

Let $\omega_0 \in \Omega$ be a given element, called compare-element, and $a_0 = a(\omega_0)$, called vision-angle.

The triplet (a_U, a_C, a_L) , where $a_U = a^{-1}(\lceil a_0 \rceil)$, $a_C = a^{-1}(\{a_0\})$, and $a_L = a^{-1}([0, a_0[)$, is called contradiction set \hat{a} .

We define:

$$\neg \hat{a} = \hat{\bar{a}} \quad \text{s.t.} \quad \bar{a}(\omega) = 1 - a(\omega) \quad \text{complementation}$$

$$\hat{a} \cup \hat{b} = \hat{v} \quad \text{s.t.} \quad v(\omega) = \begin{cases} a(\omega) \vee b(\omega) & \omega \in a_U \cup b_U \\ a(\omega) \wedge b(\omega) & \text{iff } \omega \in a_L \cap b_L \\ a_0 \wedge b_0 & \text{otherwise} \end{cases} \quad \text{union}$$

$$\hat{a} \cap \hat{b} = \hat{w} \quad \text{s.t.} \quad w(\omega) = \begin{cases} a(\omega) \wedge b(\omega) & \omega \in a_U \cap b_U \\ a(\omega) \vee b(\omega) & \text{iff } \omega \in a_L \cup b_L \\ a_0 \vee b_0 & \text{otherwise} \end{cases} \quad \text{intersection}$$

It results:

$$\neg \hat{a} = (a_L, a_C, a_U)$$

$$\hat{a} \cup \hat{b} = (a_U \cup b_U, (a_C \cap b_C) \cup (a_C \cap b_L) \cup (a_L \cap b_C), a_L \cap b_L)$$

$$\hat{a} \cap \hat{b} = (a_U \cap b_U, (a_C \cap b_C) \cup (a_C \cap b_U) \cup (a_U \cap b_C), a_L \cup b_L)$$

$$\neg (\hat{a} \cup \hat{b}) = \neg \hat{a} \cap \neg \hat{b}$$

$$\neg (\hat{a} \cap \hat{b}) = \neg \hat{a} \cup \neg \hat{b} .$$

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