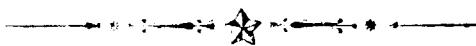


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Fuzzy permutation with lower solution of Fuzzy relation equation

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## Abstract

The lower solutions of Fuzzy relation equation are discussed; The reflexivable permutation  $E$ , operator  $\hat{\tau}$  and the operation of  $E$  with  $\hat{\tau}$  are defined in this paper. Fuzzy permutation  $p$  represented in [1], [2], it is to find the lower solutions of Fuzzy relation equation and their necessary and sufficient condition, thus a complete method for seeking lower solutions is established. If this method is tabulated, it is convenience to find all lower solutions without repetition.

## A Reflexivable permutation

Let  $U=\{1, 2, \dots, n\}$ ,  $V=\{1, 2, \dots, m\}$ . A mapping  $p$ , corresponding each point of  $U$  to a element of  $\mathcal{P}(V)$  (Fuzzy power set), is called a Fuzzy permutation and denoted by

$$p = \begin{pmatrix} 1 & 2 & \cdots & n \\ \underline{R_1} & \underline{R_2} & \cdots & \underline{R_n} \end{pmatrix}$$

where  $\underline{R_i} = \underline{r_{i1}} + \underline{r_{i2}} + \cdots + \underline{r_{im}} \in \mathcal{P}(V)$ ,  $i \in U$ . For convenience, we rewrite  $\underline{R_i}$  as the form

$$\underline{R_i} = \left( \frac{\underline{r_{i1}}}{1}, \frac{\underline{r_{i2}}}{2}, \dots, \frac{\underline{r_{im}}}{m} \right)$$

If  $U=V$  and  $\underline{R_i}$  is a Fuzzy number, then  $p$  is called a regular Fuzzy permutation.  $\underline{R_i}$  is also called the Fuzzy image of  $i$  under Fuzzy permutation  $p$ , which is denoted as  $p(i)=\underline{R_i}$ . Furthermore, we take  $p(i_1, i_2), \dots, p(V)$  as  $\{\underline{R_{i1}}, \underline{R_{i2}}, \dots, \underline{R_n}\}$  respectively.  $R_{i(j)}$  denote the degree of point  $i$  permuting by point  $j$ .

Let  $t_{ji}=r_{ij}$ ,  $\underline{T_j} = \left( \frac{t_{j1}}{1}, \frac{t_{j2}}{2}, \dots, \frac{t_{jn}}{n} \right) \in \mathcal{P}(U)$ , then  $p' = \begin{pmatrix} 1 & 2 & \cdots & n \\ \underline{T_1} & \underline{T_2} & \cdots & \underline{T_n} \end{pmatrix}$  is called the transposed permutation.

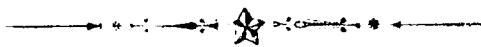
For the multiplication operation "o" defined in [1], we have generally  $p \cdot p' \neq p' \cdot p$ .

Definition 1 Let  $\underline{R_i} = e_{ii}/i$ ,  $e_{ii} \in [0, 1]$

$$E = \begin{pmatrix} 1 & 2 & \cdots & n \\ (\underline{e_1}) & (\underline{e_{12}}) & \cdots & (\underline{e_{1n}}) \end{pmatrix}$$

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is called a Fuzzy reflexivable permutation, or, reflexivable permutation.

When  $\forall i \in U, e_{ii} = 1$ ,  $E$  is the complete reflexivable permutation, i.e. the unit permutation.

When  $\forall i \in U, e_{ii} = 0$ ,  $E$  is the inverse reflexivable permutation.

Under the finite universe of discourse (or universe), Fuzzy relation is represented by the Fuzzy matrix, Fuzzy graph and Fuzzy permutation, therefore, we can establish their corresponding relation.

Example 1 Let  $R = \begin{pmatrix} 0.2 & 0.7 & 0.6 & 0.3 \\ 0.7 & 0 & 0.6 & 0.4 \\ 0.6 & 0.6 & 0 & 0.1 \\ 0.3 & 0.4 & 0.1 & 0 \end{pmatrix}$  be Fuzzy relation. The corresponding relation of the Fuzzy matrix, Fuzzy graph and the Fuzzy permutation is

$$P = \begin{pmatrix} 1 & 2 & 3 & 4 \\ (\frac{0.2}{3}, \frac{0.7}{4}) & (\frac{0.3}{2}, \frac{0.6}{1}) & (\frac{0.1}{3}, \frac{0.6}{2}) & (\frac{0.4}{1}) \end{pmatrix}$$

$$A = \begin{pmatrix} 0 & 0.6 & 0 & 0.4 \\ 0 & 0.3 & 0.6 & 0 \\ 0.2 & 0 & 0.1 & 0 \\ 0.7 & 0 & 0 & 0 \end{pmatrix} \Leftrightarrow \begin{array}{c} \text{Fuzzy Graph: } \\ \text{Vertices: } 1, 2, 3, 4 \\ \text{Edges: } (1, 2) \text{ with weight } 0.7, (1, 3) \text{ with weight } 0.6, (1, 4) \text{ with weight } 0.3, \\ (2, 1) \text{ with weight } 0.6, (2, 3) \text{ with weight } 0.2, (2, 4) \text{ with weight } 0.4, \\ (3, 1) \text{ with weight } 0.6, (3, 2) \text{ with weight } 0.1, (3, 4) \text{ with weight } 0.4, \\ (4, 1) \text{ with weight } 0.4, (4, 2) \text{ with weight } 0.3 \end{array} \quad (\text{fig. 1})$$

The difference of general permutation and Fuzzy permutation is that the general permutation corresponds point to point, where Fuzzy permutation corresponds point to Fuzzy Set.

In example 1,  $P$  replaces point "1" by  $R_1 = (\frac{0.2}{3}, \frac{0.7}{4})$ , points "3", "4" in  $\text{Supp}(R_1) = \{3, 4\}$  by  $R_3 = (\frac{0.1}{3}, \frac{0.6}{2})$  and  $R_4 = (\frac{0.4}{1})$  respectively, points "1", "2" and "3" in  $\text{Supp}(R_3) \cup \text{Supp}(R_4)$  by  $R_1 = (\frac{0.2}{3}, \frac{0.7}{4})$ ,  $R_2 = (\frac{0.3}{2}, \frac{0.6}{1})$  and  $R_3 = (\frac{0.1}{3}, \frac{0.6}{2})$  respectively, and so on.

B Characteristic Fuzzy set and  $\hat{\delta}$  operator

In [2], The Fuzzy relation equation  $A \cdot Z = B$  has been transform into iso-solution equation  $A^* \cdot Z = B$ . Where  $A^*$  is iso-solution matrix of  $A$  and  $g_i^* \in \{0, b_j\}$ .

Example 2 The iso-solution matrix of Fuzzy relation equation

$$\begin{pmatrix} 0.8 & 0.9 & 0.7 & 0.5 & 0.4 & 0.1 \\ 0.8 & 0.7 & 0.8 & 0.7 & 0.6 & 0.6 \\ 0.3 & 0.8 & 1 & 0.9 & 0.6 & 0.9 \\ 0.6 & 0.6 & 0.7 & 0.8 & 0.8 & 0.8 \\ 0.4 & 0.3 & 0.5 & 0.4 & 0.2 & 0.6 \\ 0.5 & 0.2 & 0.4 & 0.3 & 0.6 & 0.5 \\ 0.2 & 0.4 & 0.3 & 0.3 & 0.5 & 0.3 \\ 0.2 & 0.2 & 0.2 & 0.1 & 0.4 & 0.4 \\ 0.3 & 0.1 & 0 & 0.3 & 0.1 & 0.3 \end{pmatrix} \circ \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \end{pmatrix} = \begin{pmatrix} 0.8 \\ 0.8 \\ 0.8 \\ 0.7 \\ 0.5 \\ 0.5 \\ 0.4 \\ 0.4 \\ 0.3 \end{pmatrix} \quad (1)$$

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$$A^* = \begin{pmatrix} 0.8 & 0.8 & 0 & 0 & 0 & 0 \\ 0.8 & 0 & 0.8 & 0 & 0 & 0 \\ 0 & 0.8 & 0.8 & 0 & 0 & 0 \\ 0 & 0 & 0.7 & 0.7 & 0 & 0 \\ 0 & 0 & 0.5 & 0 & 0 & 0.5 \\ 0.5 & 0 & 0 & 0 & 0 & 0.5 \\ 0 & 0.4 & 0 & 0 & 0.4 & 0 \\ 0 & 0 & 0 & 0 & 0.4 & 0.4 \\ 0.3 & 0 & 0 & 0.3 & 0 & 0.3 \end{pmatrix} \quad (2)$$

In the rest of the paper, we suppose  
 $b_1 \geq b_2 \geq \dots \geq b_m$

and  $b_1 = b_2 = \dots = b_t = \lambda_1$ ,  $b_{t+1} = \dots = b_r = \lambda_2$ , ...,  $b_{r-1} = \dots = b_m = b_{t_k} = \lambda_t$

whenever we discuss iso-solution equation.

The element  $a_{ij}^*$  of iso-solution matrix  $A^*$  is equal to  $b_j$  or greater to  $b_j$ .  $\underline{R}_i$  is Fuzzy set which is consisted of all non-zero element of  $i$ -th column of  $A^*$ . Obviously, Fuzzy image of  $i$  regarded as Fuzzy vector.

**Definition 2** Let  $A^*$  corresponding to Fuzzy permutation  $p$ , and  $\lambda \in [0, 1]$ ,  $\exists j \in V$ . such that

$$\underline{R}_i(j) = \lambda$$

$$\bigcup_{i \in I} \{ \underline{R}_{i(j)} \mid \underline{R}_{i(j)}(j) = \lambda \}$$

is called the lower bound Fuzzy set of  $U$  for  $\lambda$ , in brief,  $\lambda$ -lower Fuzzy set, and

$$I = \{i_1, i_2, \dots, i_k\} \subset U$$

is called the  $\lambda$ -lower index set

Obvious,  $(\underline{R}_{i_1} \cup \underline{R}_{i_2} \cup \dots \cup \underline{R}_{i_k})_\lambda = \{j_1, j_2, \dots, j_r\} = J \subset V$ ,  $J$  is called the  $\lambda$ -lower cuts set.

In example 2, The  $A^*$  corresponds to  $p$ :

$$p = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ (\frac{0.8}{1}, \frac{0.8}{2}, \frac{0.5}{6}, \frac{0.3}{9}) & (\frac{0.8}{1}, \frac{0.8}{3}, \frac{0.4}{7}) & (\frac{0.8}{2}, \frac{0.8}{3}, \frac{0.7}{4}, \frac{0.5}{5}) & (\frac{0.7}{4}, \frac{0.3}{5}) & (\frac{0.4}{7}, \frac{0.4}{8}) & (\frac{0.5}{5}, \frac{0.5}{6}, \frac{0.4}{8}, \frac{0.3}{7}) \end{pmatrix}$$

Let  $\lambda = 0.8$ , there exist  $j=1, 2, 3 \in V$ , such that

$$\underline{R}_1(1) = \underline{R}_2(3) = \underline{R}_3(2) = 0.8$$

so  $(\frac{0.8}{1}, \frac{0.8}{2}, \frac{0.5}{6}, \frac{0.3}{9}) \cup (\frac{0.8}{1}, \frac{0.8}{3}, \frac{0.4}{7}) \cup (\frac{0.8}{2}, \frac{0.8}{3}, \frac{0.7}{4}, \frac{0.5}{5})$  is  $\lambda$ -lower bound Fuzzy set, and  $I = \{1, 2, 3\} \subset U$  is  $\lambda$ -lower bound index set.

Because  $(\underline{R}_1 \cup \underline{R}_2 \cup \underline{R}_3)_{0.8} = \{1, 2, 3\} \subset V = \{1, 2, \dots, 9\}$ , so  $\lambda$ -lower bound cuts is  $J = \{1, 2, 3\}$ .

**Definition 3** Suppose that there is  $\lambda$ -lower bound index set  $I = \{i_1, i_2, \dots, i_k\}$ , if  $\exists I_t \subset I$  ( $I_t = \{i_{t_1}, i_{t_2}, \dots, i_{t_r}\}$ ), such that

$$(\underline{R}_{i_1} \cup \underline{R}_{i_2} \cup \dots \cup \underline{R}_{i_k})_\lambda = J$$

but  $\forall I_s \subseteq I_t$  ( $s \geq 1$ ),  $I_s = \{i_{s_1}, i_{s_2}, \dots, i_{s_s}\}$ , then

$$(\underline{R}_{i_{s_1}} \cup \underline{R}_{i_{s_2}} \cup \dots \cup \underline{R}_{i_{s_s}}) \nsubseteq J$$

thus  $(\underline{R}_{i_1} \cup \underline{R}_{i_2} \cup \dots \cup \underline{R}_{i_{t_r}})$  is called the characteristic Fuzzy set, and  $I_t$  is called the characteristic index set. Without any confusion the characteristic index set be represented by  $\{i_1, i_2, \dots, i_t\}$  and characteristic Fuzzy set by  $\underline{R}_{i_1} \cup \underline{R}_{i_2} \cup \dots \cup \underline{R}_{i_t}$ .

Obviously, the characteristic index set is chosen by means of more than one ways in  $I$ , and when  $t_1 \neq t_2$ , then  $I_{t_1} \notin I_{t_2}$  or  $I_{t_2} \notin I_{t_1}$ , denote  $I_t^c$  by  $I \setminus I_t$ .

**Definition 4** Let Characteristic Fuzzy Set be  $\underline{R}_{i_1} \cup \underline{R}_{i_2} \cup \dots \cup \underline{R}_{i_k}$  and  $\lambda \in [0, 1]$ . Defining reflexivable permutation of  $p$  is

$$E = \begin{pmatrix} 1 & 2 & \dots & m \\ \underline{E}_1 & \underline{E}_2 & \dots & \underline{E}_m \end{pmatrix}, \quad \underline{E}_j = \frac{\underline{e}_{ij}}{j}, \quad j \in V$$

for characteristic Fuzzy set, where

$$\underline{e}_{ij} = \begin{cases} 0, & j \text{ satisfying } 0 \leq \underline{R}_i(j) < \lambda \quad (i \in I_t) \\ \lambda, & \text{other } j. \end{cases}$$

For example, we take out characteristic index set  $I_t = \{1, 2\}$  in  $p$ , thus

$$E_{I_t} = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ (\frac{0.8}{1}) & (\frac{0.8}{2}) & (\frac{0.8}{3}) & (\frac{0.8}{4}) & (\frac{0.8}{5}) & (\frac{0}{6}) & (\frac{0}{7}) & (\frac{0.8}{8}) & (\frac{0}{9}) \end{pmatrix}$$

**Definition 5** Let reflexivable permutation of Fuzzy permutation  $p$  for characteristic Fuzzy set be  $E$ , we definite the operation  $\hat{\delta}$ :

$$p' = E \hat{\delta} p = \begin{pmatrix} 1 & 2 & \dots & m \\ \underline{E}_1 & \underline{E}_2 & \dots & \underline{E}_m \end{pmatrix} \hat{\delta} \begin{pmatrix} 1 & 2 & \dots & n \\ \underline{R}_1 & \underline{R}_2 & \dots & \underline{R}_n \end{pmatrix} = \begin{pmatrix} 1 & 2 & \dots & n \\ \underline{G}_1 & \underline{G}_2 & \dots & \underline{G}_n \end{pmatrix}$$

where  $\underline{G}_i = \left( \frac{g_{i1}}{1}, \frac{g_{i2}}{2}, \dots, \frac{g_{im}}{m} \right)$ ,  $g_{ij} = \begin{cases} \underline{e}_{ij} \hat{\delta} \underline{r}_{ij}, & (i \in I_t) \\ \underline{e}_{ij} \hat{\delta} \underline{r}_{ij}, & (\text{otherwise}). \end{cases}$

and  $\hat{\wedge}$ ,  $\hat{\delta}$  are definited respectively by

$$a \hat{\wedge} b \triangleq \begin{cases} b, & (a \geq b) \\ 0, & (a < b), \end{cases} \quad a \hat{\delta} b \triangleq \begin{cases} b, & (a > b) \\ 0, & (a \leq b). \end{cases}$$

For all  $\lambda \in [0, 1]$ , the definitions 4 and definition 5 repeatedly apply to  $p$ , and let  $\underline{R}_i^{(1)}(j) = x_i$ , then

$$\underline{Z} = \begin{pmatrix} \underline{X}_1 \\ \underline{X}_2 \\ \vdots \\ \underline{X}_n \end{pmatrix} = \begin{pmatrix} \underline{R}_1^{(1)}(j) \\ \underline{R}_2^{(1)}(j) \\ \vdots \\ \underline{R}_n^{(1)}(j) \end{pmatrix} \quad (3)$$

is a lower solution of Fuzzy relation equation. In example 2,  $(0.8, 0.8, 0.7, 0, 0, 0.5)^T$  is a lower solution

### C Necessary and sufficient condition of Fuzzy relation equation have a lower solution

**Theorem 1** Suppose that  $A^* \cdot \underline{Z} = B$ ,  $A^* \rightarrow p$ . Fuzzy relation equation has solution iff

$$\text{Supp}\{\underline{R}_1, \underline{R}_2, \dots, \underline{R}_n\} = \text{Supp}[p(U)] = V.$$

**Proof** Because  $A^* \Leftrightarrow p$ , the  $\underline{R}_1, \underline{R}_2, \dots, \underline{R}_n$  correspond respectively to 1-th, 2-th, ..., n-th column matrix of  $A^*$ . Obviously, row order number set of  $A^*$  is  $V$ , Column order number set of  $A^*$  is  $U$ .

**Necessity** Because  $A^* \circ \underline{x} = B$  has solution, then

$$\begin{cases} \bigvee_{i=1}^n (a_{ji}^* \wedge x_i) \leq b_j & (\forall j \in V), \\ (a_{j,i_0}^* \wedge x_{i_0}) = b_j & (\forall j \in V, \exists i_0 \in U). \end{cases}$$

So that, there exists non-zero element  $a_{ji}^*$  of  $A^*$  at least in each row of  $A^*$ ; and because  $a_{ji}^* = r_{ij} (= p_i(j))$ , i.e. non-zero  $r_{ij}$ . then there exist non-zero element  $r_{ij}$  in  $\underline{R}_{i_0}$  or  $\underline{R}_{i_1}, \dots, \underline{R}_{i_m}$ , i.e. there are non-zero elements in

$$\underline{R}_1 \cup \underline{R}_2 \cup \dots \cup \underline{R}_n.$$

and  $j$  is any element of  $V$ , so that

$$\text{Supp}\{\underline{R}_1 \cup \underline{R}_2 \cup \dots \cup \underline{R}_n\} = \text{Supp}(p(U)) = V$$

**Sufficiency** Suppose that  $A^* \circ \underline{x} = B$ ,  $A^* \Leftrightarrow p$ , then  $a_{ji}^* \in \{0, b_j\}$ .

Let  $\text{Supp}(p(U)) = V$ . then there are  $m$  non-zero  $r_{ij}$  in  $\underline{R}_1 \cup \underline{R}_2 \cup \dots \cup \underline{R}_n$  at least.

If  $a_{1i_1}^* = b_1, a_{2i_2}^* = b_2, \dots, a_{mi_m}^* = b_m$ ,

obviously  $(a_{j,i_1}^* \wedge x_1) \vee (a_{j,i_2}^* \wedge x_2) \vee \dots \vee (a_{j,i_m}^* \wedge x_m) \leq b_j \quad (j=1, 2, \dots, m)$

therefore, for each  $j$ , there is at least one element of  $I$ , say  $i_0$ , then

$$a_{j,i_0}^* \wedge x_{i_0} = b_j$$

thus

$$A^* \circ \underline{x} = B.$$

**Corollary 1**  $\text{Supp}[p^*(j)] \neq \emptyset$

**Definition 6** Suppose that  $A^* \circ \underline{x} = B$ , for each row (column) of  $A^*$  operated "V" according to Column (row)  $a_{ji}^*$ , then  $A^*$  becomes the column (row) matrix, its called column (row) projection. It denoted by

$$\bigvee_{i \in U} A^* \quad (\bigvee_{j \in V} A^*)$$

**Theorem 2** Suppose that there is a solution in  $A^* \circ \underline{x} = B$ , then

$$1^\circ \bigvee_{i \in U} A^* = B.$$

$$2^\circ A^* \circ (\bigvee_{j \in V} A^*)^T = B \quad (T \text{ representing transpose operation}).$$

**Proof** 1° Because  $A^* \circ \underline{x} = B$  has solution, thus by the result of corollary 1 of theorem 1,  $\forall j \in V$ , we obtain

$$\text{Supp}[p^*(j)] \neq \emptyset \quad (\text{each row nonempty set of } A^*).$$

Because  $a_{ji}^* \in \{0, b_j\}$ , thus, for any  $j$  such that

$$\max\{a_{j1}^*, a_{j2}^*, \dots, a_{jn}^*\} = b_j.$$

Because  $j$  is any element, then

$$\bigvee_{i \in U} A^* = B.$$

2° If  $A^* \circ \underline{x} = B$  has a solution  $\underline{x}$ ,

$$\begin{cases} (a_{11}^* \wedge x_1) \vee \dots \vee (a_{1i}^* \wedge x_i) \vee \dots \vee (a_{1n}^* \wedge x_n) = b_1, \\ (a_{21}^* \wedge x_1) \vee \dots \vee (a_{2i}^* \wedge x_i) \vee \dots \vee (a_{2n}^* \wedge x_n) = b_2, \\ \dots \quad \dots \quad \dots \quad \dots \quad \dots \quad \dots \quad \dots \\ (a_{m1}^* \wedge x_1) \vee \dots \vee (a_{mi}^* \wedge x_i) \vee \dots \vee (a_{mn}^* \wedge x_n) = b_m \end{cases} \quad (4)$$

thus

$$\begin{cases} a_{ji}^* \wedge x_i \leq b_j, & (\forall j \in V, \forall i \in U) \\ a_{j_{i_0}}^* \wedge x_{i_0} = b_j, & (\forall j \in V, \exists i_0 \in U) \end{cases}$$

Obviously,  $a_{j_{i_0}}^* \geq b_j$ ,  $x_{i_0} \geq b_j$  and  $a_{ji} \in \{0, b_j\}$ , thus  $a_{j_{i_0}}^* = b_j$ , then  $\forall j \in V, \exists i_0 \in U$ , such that  $x_{i_0} \geq b_j = a_{j_{i_0}}^*$ .

$$\text{Let } x'_i = \max\{a_{1i}^*, a_{2i}^*, \dots, a_{ni}^*\} \quad (\forall j \in V, \exists i \in U)$$

Because  $\text{supp}(p^*(j)) \neq \emptyset$  ( $\forall j \in V$ ), for  $i$ -th column there is no harm in letting  $a_{1i}^* = b_1$ .

$$\text{therefore } x'_i \geq \max\{a_{1i}^*, a_{2i}^*, \dots, a_{ni}^*\} = a_{1i}^* = b_1.$$

thus we replaces  $x_{i_1}$  of first equation of (4) by  $x'_{i_1}$ , then

$$(a_{11}^* \wedge x_1) \vee \dots \vee (a_{i_1}^* \wedge x'_{i_1}) \vee \dots \vee (a_{nn}^* \wedge x_n) = (a_{11}^* \wedge x_1) \vee \dots \vee (b_1 \wedge x'_{i_1}) \vee \dots \vee (a_{nn}^* \wedge x_n) = b_1.$$

Similarly, for  $i_2$ , following formula holds

$$x'_{i_2} \geq \max\{a_{1i_2}^*, a_{2i_2}^*, \dots, a_{ni_2}^*\} = a_{2i_2}^* = b_2 \leq b_1,$$

then we replace  $x_{i_1}$  together with  $x_{i_2}$  by  $x'_{i_1}$ ,  $x'_{i_2}$  respectively in the first two equations of (4), therefore

$$(a_{11}^* \wedge x_1) \vee \dots \vee (a_{i_1}^* \wedge x'_{i_1}) \vee \dots \vee (a_{i_2}^* \wedge x'_{i_2}) \vee \dots \vee (a_{nn}^* \wedge x_n) = b_1,$$

$$(a_{11}^* \wedge x_1) \vee \dots \vee (a_{i_1}^* \wedge x'_{i_1}) \vee \dots \vee (a_{i_2}^* \wedge x'_{i_2}) \vee \dots \vee (a_{2n}^* \wedge x_n) = b_2.$$

Similarly,  $\forall i \in U$ , we solve  $x'_1, x'_2, \dots, x'_n$ , let  $\bar{x}' = (x'_1, x'_2, \dots, x'_n)^T$ , then  $A^* \cdot \bar{x}' = B$ .

where

$$\bar{x}' = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} \max\{a_{11}^*, a_{21}^*, \dots, a_{n1}^*\} \\ \max\{a_{12}^*, a_{22}^*, \dots, a_{n2}^*\} \\ \vdots \\ \max\{a_{1n}^*, a_{2n}^*, \dots, a_{nn}^*\} \end{pmatrix} = (\bigvee_{j \in V} A^*)^T$$

$$\text{i.e. } A^* \circ (\bigvee_{j \in V} A^*)^T = B.$$

Because  $A^* \rightleftharpoons p$ , evidently

$$\begin{aligned} \bigvee_{j \in V} A^* &= (\bigvee_{j \in V} p(1), \bigvee_{j \in V} p(2), \dots, \bigvee_{j \in V} p(n)) \\ &= (\bigvee_{j \in V} R_1, \bigvee_{j \in V} R_2, \dots, \bigvee_{j \in V} R_n) \end{aligned}$$

$(\bigvee_{j \in V} p(1), \dots, \bigvee_{j \in V} p(n))$  denoted by  $\bigvee_{j \in V} p$ , then

Corollary 2 Suppose that there is a solution in  $A^* \cdot \bar{x} = B$ , then  $A^* \circ (\bigvee_{j \in V} p)^T = B$ .

For convenience, the  $p$  is resolved by the union of Fuzzy permutation, for instance

$$\begin{aligned} p &= \left( \begin{array}{ccc} 1 & 2 & 3 \\ \left(\frac{0.8}{1}, \frac{0.8}{2}, \frac{0.5}{6}, \frac{0.3}{9}\right) & \left(\frac{0.8}{1}, \frac{0.8}{3}, \frac{0.4}{7}\right) & \left(\frac{0.5}{5}, \frac{0.5}{6}, \frac{0.4}{8}\right) \end{array} \right) \\ &= \left( \begin{array}{ccc} 1 & 2 & 3 \\ \left(\frac{0.8}{1}, \frac{0.8}{2}\right) & \left(\frac{0.4}{7}\right) & \left(\frac{0.5}{5}, \frac{0.3}{6}\right) \end{array} \right) \cup \left( \begin{array}{ccc} 1 & 2 & 3 \\ \left(\frac{0.5}{6}, \frac{0.3}{9}\right) & \left(\frac{0.8}{1}, \frac{0.8}{3}\right) & \left(\frac{0.4}{8}\right) \end{array} \right) \end{aligned}$$

then, we have

Definition 7 Let  $p$  be Fuzzy permutation and

$$p = p^1 \cup p^2 \cup \dots \cup p^K, \text{ where } p^l = \left( \begin{array}{cccc} 1 & 2 & \cdots & n \\ \underbrace{R_1^l} & \underbrace{R_2^l} & \cdots & \underbrace{R_n^l} \end{array} \right) \quad (1 \leq l \leq K)$$

and  $R_i^l = R_i^1 \cup R_i^2 \cup \dots \cup R_i^l$  ( $1 \leq i \leq n$ ).  $p^1 \cup p^2 \cup \dots \cup p^K$  is called resolution of  $p$ .

By means of this representation of  $p$  we obtain

$$\begin{aligned} A^* \circ (\bigvee_{j \in V} p)^T &= A^* \circ (\bigvee_{j \in V} (p^1 \cup p^2 \cup \dots \cup p^k))^T \\ &= A^* \circ (\bigvee_{j \in V} p^1)^T \cup A^* \circ (\bigvee_{j \in V} p^2)^T \cup \dots \cup A^* \circ (\bigvee_{j \in V} p^k)^T = B, \end{aligned}$$

evidently

$$A^* \circ (\bigvee_{j \in V} p^k)^T \subset B \quad (k \geq 1)$$

According to definition 7 we obtain

Theorem 3 If  $A^* \circ (\bigvee_{j \in V} p)^T = B$ , then there exists certain  $p^t$  ( $p^t \subseteq p$ ) such that

$$A^* \circ (\bigvee_{j \in V} p^t)^T = B \quad (1 \leq t \leq k)$$

Definition 8 If  $A^* \circ (\bigvee_{j \in V} p^t)^T = B$ ,  $\forall p^{t_i} \not\subseteq p^t$ ,  $A^* \circ (\bigvee_{j \in V} p^{t_i})^T \not\subseteq B$ , then  $(\bigvee_{j \in V} p^t)^T$  is called lower solution of  $A^* \circ B = B$ .

Theorem 4 Let  $\mathcal{Z}_0$  be lower solution of  $A^* \circ B = B$ , then  $\exists x_i \in \mathcal{Z}_0$ , such that  
 $x_i = \max\{b_1, b_2, \dots, b_m\}$

where  $x_i$  is component of  $\mathcal{Z}_0$ .

Proof If  $(\bigvee_{j \in V} p^t)^T$  is lower solution of  $A^* \circ \mathcal{Z} = B$ , then

$$A^* \circ (\bigvee_{j \in V} p^t)^T = B$$

thus  $(\bigvee_{j \in V} p^t) = (\bigvee_{j \in V} R_1^t, \bigvee_{j \in V} R_2^t, \dots, \bigvee_{j \in V} R_n^t) = (x_1, x_2, \dots, x_n)$

there is no harm in that let maximum component of  $(x_1, x_2, \dots, x_n)$  be  $x_i$ . Certainly there is

$$\begin{aligned} \bigvee_{j \in V} R_i^t &= \max\{a_{1i}^t, a_{2i}^t, \dots, a_{ni}^t\} \quad (i \in U) \\ &= x_i = b_i \end{aligned}$$

This theorem represents that there must exist maximum component at least in lower solution, the maximum component is  $b_i$ .

For convenience, we lead into following sign:

(1) Maximum component of solution of  $A^* \circ B = B$  denoted by « $\mathcal{Z}$ », for instance, « $(0.8, 0.8, 0.7, 0, 0.4, 0)$ »  $\cong (0.8, 0.8, 0, 0, 0, 0)$

(2) Let  $x_{i_1}, x_{i_2}, \dots, x_{i_k}$  be  $k$  components of  $\mathcal{Z}$ , we have

$$\{\{x_{i_k}\}\} \cong (0, \dots, 0, x_{i_1}, 0, \dots, 0, x_{i_2}, 0, \dots, 0, x_{i_k}, 0, \dots, 0)$$

for instance,  $\{0.8, 0.4\} = (0.8, 0, 0, 0, 0.4, 0)$ .

Definition 9 If  $A^* \circ \mathcal{Z} = B$ , and  $b_{j_1} = b_{j_2} = \dots = b_{j_t} = \lambda_t$ , we take out

$$\left\{ \begin{array}{l} (a_{j_1} \wedge x_1) \vee (a_{j_2} \wedge x_2) \vee \dots \vee (a_{j_n} \wedge x_n) = b_{j_1}, \\ (a_{j_2} \wedge x_1) \vee (a_{j_3} \wedge x_2) \vee \dots \vee (a_{j_n} \wedge x_n) = b_{j_2}, \\ \dots \\ (a_{j_t} \wedge x_1) \vee (a_{j_{t+1}} \wedge x_2) \vee \dots \vee (a_{j_n} \wedge x_n) = b_{j_t} \end{array} \right. \quad (5)$$

from  $A^* \circ \mathcal{Z} = B$ . (5) is called the  $\lambda_t$ -level lumped Fuzzy relation equation, and denoted by  $A_{J(\lambda_t)}^* \circ \mathcal{Z} = B_{J(\lambda_t)}$ , in brief,  $\lambda_t$ -lumped equation. Evidently, we have

Theorem 5 The solutions of  $A^* \circ \mathcal{Z} = B$  is also solutions of  $\lambda$ -lumped equation.

Theorem 6 If  $A^* \circ \mathcal{Z} = B$  has solutions, maximum component of lower solution are  $x_{i_1}, x_{i_2}, \dots, x_{i_k}$  iff  $I_t = \{i_1, i_2, \dots, i_k\}$  is characteristic index set.

Proof. Let maximum component of lower solution of  $A^* \cdot \underline{x} = B$  be  $x_i = x_{i_1} = \dots = x_{i_k} = \lambda_1$ . Thus, we can take one  $\lambda_1$ -lumped equation with  $j$  rows from  $A^* \cdot \underline{x} = B$ .

$$A_{J(\lambda_1)}^* \cdot \underline{x} = B_{J(\lambda_1)}$$

Where  $J(\lambda_1) = \{1, 2, \dots, j\}$  represents the set defined by  $\lambda_1$ . Evidently,  $A_{J(\lambda_1)}^*$  is a  $j \times n$  matrix, and  $A_{J(\lambda_1)}^*$  correspond to  $p_{J(\lambda_1)}$ , lower solution of  $A^* \cdot \underline{x} = B$  is also lower solution of  $A_{J(\lambda_1)}^* \cdot \underline{x} = B_{J(\lambda_1)}$ . Then, there are solutions of  $A_{J(\lambda_1)}^* \cdot \underline{x} = B_{J(\lambda_1)}$ , and by theorem 1,

$$\text{Supp}[p_{J(\lambda_1)}(U)] = J(\lambda_1) = \{1, 2, \dots, j\}$$

holds. There is no harm in letting

$$\bigvee_{j \in U} p_{J(\lambda_1)}(i_j) = \lambda_1 \quad (i \in I = \{i_1, i_2, \dots, i_r\} \subset U),$$

$$\bigvee_{j \in U} p_{J(\lambda_1)}(i'_j) < \lambda_1 \quad (i' \notin I).$$

Then,

$$\text{Supp}[p_{J(\lambda_1)}(I)] = J$$

holds. Evidently

$$(p_{J(\lambda_1)}(i_1), p_{J(\lambda_1)}(i_2), \dots, p_{J(\lambda_1)}(i_r))_{\lambda_1} = J$$

$x_{i_1}, x_{i_2}, \dots, x_{i_k}$  are maximum components of lower solution of  $A^* \cdot \underline{x} = B$  ( $x_{i_1}, x_{i_2}, \dots, x_{i_k}$  was known).

Thus for  $\{i_1, i_2, \dots, i_k\} \subset \{i_1, i_2, \dots, i_r\}$  ( $k \leq r$ ), then

$$A_{J(\lambda_1)}^* \cdot (\{\bigvee_{j \in J(\lambda_1)} p_{J(\lambda_1)}(i_1, i_2, \dots, i_k)\})^T = B_{J(\lambda_1)}.$$

If  $\{i'_1, i'_2, \dots, i'_s\} \not\subseteq \{i_1, i_2, \dots, i_k\}$ , then

$$A_{J(\lambda_1)}^* \cdot (\{\bigvee_{j \in J(\lambda_1)} p_{J(\lambda_1)}(i'_1, i'_2, \dots, i'_s)\})^T \not\subseteq B_{J(\lambda_1)}.$$

thus,  $I_t = \{i_1, i_2, \dots, i_k\}$  is characteristic index set by means of definition 3.

Vice versa.

**Corollary 3** Lower solution of  $\lambda_1$ -lumped equation is also  $(\{\bigvee_{i \in I_t} p(i)\})^T$ , ( $i \in I_t$ ) where  $I_t$  is a characteristic index set.

**Definition 10** If lower solution of  $\lambda_1$ -lumped equation is  $(\{\bigvee_{j \in J} p_{J(\lambda_1)}(i_j)\}_{i \in I_t})^T$ , and there exists  $j' \in \text{Supp}(p(i_1), p(i_2), \dots, p(i_k)) = \text{Supp}(p(I_t))$  but  $j' \notin J$ . Let  $J'$  be the set of all elements  $j'$  in  $\text{Supp}(p(I_t))$  and let  $J_t = J \cup J'$ .  $A_{J_t(\lambda_1)}^* \cdot \underline{x} = B_{J_t(\lambda_1)}$  is called the  $\lambda_1$ -lumped extension equation of  $A_{J(\lambda_1)}^* \cdot \underline{x} = B_{J(\lambda_1)}$

**Theorem 7** Lower solution of  $\lambda_1$ -lumped equation is a lower solution of  $\lambda_1$ -lumped extension equation. Vice versa.

Proof.  $p_{J_t(\lambda_1)}^* \leftarrow A_{J_t(\lambda_1)}^*$ , evidently  $\bigvee_{j \in J} p_{J(\lambda_1)}(i_1, i_2, \dots, i_k) = \bigvee_{j \in J_t} p_{J_t(\lambda_1)}(i_1, i_2, \dots, i_k)$ , because,

$$A_{J(\lambda_1)}^* \cdot (\bigvee_{j \in J} p_{J(\lambda_1)}(i_1, i_2, \dots, i_k))^T = B_{J(\lambda_1)}$$

then

$$A_{J_t(\lambda_1)}^* \cdot (\bigvee_{j \in J_t} p_{J_t(\lambda_1)}(i_1, i_2, \dots, i_k))^T = B_{J_t(\lambda_1)}.$$

Therefore

$$I_t = \{i_1, i_2, \dots, i_k\}$$
 is characteristic index set.

The theorem 7 presents an important fact: If we seek out lower solution of  $\lambda_1$ -lumped equation, then this solution satisfies voluntarily for  $\lambda_1$ -lumped extension equation. Therefore after determining characteristic Fuzzy set for  $\lambda_1$  in  $p$ . We discard the intersection " $\cap$ " of this characteristic Fuzzy set and non-characteristic Fuzzy set from  $p(U \setminus I_t)$ , i.e.  $p(U \setminus I_t) \setminus p(I_t) \cap p(U \setminus I_t)$ , afterward operate  $\bigvee_{j \in J} p_{J(\lambda_1)}(i_1, i_2, \dots, i_k)$ .

**Definition 11** If solution of  $A^* \circ Z = B$  exists, thus we discard  $\lambda_i$ -lumped extension equation from  $A^* \circ Z = B$ , therefore the remaining part of  $A^* \circ Z = B$  is called  $\lambda_i$ -surplus lumped equation.

**Theorem 8** The second-class component are  $X_{i_k+1}, X_{i_k+2}, \dots, X_{i_k+k_2}$  iff  $I_e^{(1)} = \{i_{k_1}+1, i_{k_1}+2, \dots, i_{k_1}+k_2\}$  is characteristic index set which is determined by the maximum level  $\lambda_i$  of  $\lambda_i$ -surplus lumped equation (it is called second-class characteristic index set)

**Proof.** The same as that of theorem 6.

Generally, we have  $t$ -th-class component and  $t$ -th-class characteristic index set, then

**Theorem 9**  $Z_0$  is lower solution of  $A^* \circ Z = B$  iff  $t$ -th-class components of  $Z_0$  correspond to  $t$ -th-class characteristic index set.

The process of determining each class component of  $A^* \circ Z = B$  according to each level  $\lambda_t$  of  $A^* \circ Z = B$ , one by one, is called the characteristic sum decomposition, denoted by

$$\lambda_1 + \lambda'_1 + \dots + \lambda'_t \quad (t \geq 1)$$

**Remark.** because  $\lambda$ -lumped extension equation is grown in  $A^* \circ Z = B$ , then, the number of  $\lambda_t$  ( $t \geq 1$ ) in characteristic sum is not equal to that of  $\lambda_t$  in  $A^* \circ Z = B$ .

To the characteristic sum decomposition, we have

**Theorem 10** There is a solution of  $A^* \circ Z = B$ ,  $E, E^{(1)}, E^{(2)}, \dots, E^{(t)}$  are fuzzy reflexivable permutation which are determined by  $t+1$ -class characteristic index set one by one, let  $E \hat{\wedge} p = p^{(1)}, E^{(1)} \hat{\wedge} p^{(1)} = p^{(2)}, \dots$ , then

$$1^\circ \quad p \supseteq E \hat{\wedge} p \supseteq E^{(1)} \hat{\wedge} p^{(1)} \supseteq \dots \supseteq p^{(t)} \hat{\wedge} E^{(t)} = p^{(t)} \neq \emptyset.$$

$$2^\circ \quad (\bigvee_{j \in V} E^{(t)} \hat{\wedge} p^{(t)})^T \text{ is lower solution of } A^* \circ Z = B$$

**Proof** We can prove 1° according to definition 4 and 5

2° First we prove  $A^* \circ (\bigvee_{j \in V} E \hat{\wedge} p)^T = B$ .

$$\begin{aligned} & \bigvee_{i \in U} [a_{ji}^* \wedge ((\bigvee_{j \in V} e_{jj} \wedge r_{ij}) \cup (\bigvee_{j \in V} e_{jj} \hat{\wedge} r_{ij}))] \\ &= \bigvee_{i \in U} [a_{ji}^* \wedge ((\bigvee_{j \in V} \lambda_j \wedge r_{ij}) \vee (\bigvee_{j \in V} \lambda_j \hat{\wedge} r_{ij}))] \\ &= \bigvee_{i \in U} [a_{ji}^* \wedge (\lambda_j \vee r_{ij}')] \end{aligned}$$

According to definition 5 we have  $r_{ij}' < \lambda_j$ ; and  $a_{ji}^* \in \{0, \lambda_j\}$ , according to theorem 1 and corollary 1, there,  $\text{supp}(p^{(t)}) \neq \emptyset$ , then  $\forall j$ , there exists at least one element in  $V$ . say  $j_0$ , such that  $a_{j_0 i}^* = b_j$ ,  $b_j = \lambda_j$ . Then

$$\bigvee_{i \in U} [a_{ji}^* \wedge (\lambda_j \vee r_{ij}')] = \bigvee_{i \in U} [a_{ji}^* \wedge \lambda_j] = \bigvee_{i \in U} [a_{ji}^* \wedge b_j] = a_{j_0 i}^* \wedge b_j = b_j.$$

Similarly,  $A^* \circ (\bigvee_{j \in V} E^{(t)} \hat{\wedge} p^{(t)})^T = B \quad (t=1, 2, \dots, t)$

therefore  $A^* \circ (\bigvee_{j \in V} p)^T = A^* \circ (\bigvee_{j \in V} E \hat{\wedge} p)^T = \dots = A^* \circ (\bigvee_{j \in V} E^{(t)} \hat{\wedge} p^{(t)})^T = A^* \circ (\bigvee_{j \in V} p^{(t)})^T = B$ .

also because 1°  $p \supseteq E \hat{\wedge} p \supseteq \dots \supseteq E^{(t)} \hat{\wedge} p^{(t)} = p^{(t)} \neq \emptyset$

thus  $\bigvee_{j \in V} p \supseteq \bigvee_{j \in V} (E \hat{\wedge} p) \supseteq \dots \supseteq \bigvee_{j \in V} p^{(t)} \neq (0, 0, \dots, 0)$ .

Then  $(\bigvee_{i=1}^n p^{(i)})^T$  is lower solution of  $A^* \cdot \mathbf{B} = \mathbf{B}$ .

$p, p^{(1)}, p^{(2)}, \dots, p^{(n)}$  are obtained by the method provide by means of theorem 10, therefore lower solution is obtained.

We illustrate the method using which the lower solution is obtained. If this method is tabulated, it is convenience to find all lower solutions without repetition. by same example 2

First, the iso-solution matrix correspond to  $p$ , i.e.  $A^* \rightarrow p$ .

Next,  $\lambda_1=0.8$  and characteristic Fuzzy set  $\{R_1, UR_2\}$  are determined. 0.8 which is supremum of  $\{R_1, UR_2\}$  is marked down in first row of table. For the functions of  $\hat{\pi}$ , we discard  $0.8/2, 0.8/3, 0.3/9, 0.4/7$  and  $0.5/6$  from  $p(U \setminus I_1)$ , i.e.  $p(3), p(4), p(5)$  and  $p(6)$ . And mark down these by "x" in first row of table. This process corresponds to the operation  $E \hat{\wedge} p$ , the selection of first-class component is completed.

The second row, we take  $\lambda_2=0.7$  and place it into second row of the table (0.7 is the supremum of characteristic Fuzzy set  $\{R_2\}$ ).  $0.7/4$  is discarded from  $R_2$  and marked by "x". This process correspond to the operation  $E'' \hat{\wedge} p''$ , then the Selection of second-class component is completed.

Take  $\lambda_3, \lambda_4$  be another number, for the same way, we can select third-class and fourth-class component.

The fifth row is sum up first to fourth row, the sixth to fourteenth row are obtained similarly, all are images of  $p^{(1)}$ , i.e. supremum of image of  $p$ . There are corresponding lower solutions of  $p^{(1)}$  image in the right side of the table, i.g. for fifth row,

$$\bigvee_{i=1}^n p^{(i)} = (0.8, 0.8, 0.5, 0.7, 0.4, 0) = X_1^T.$$

$X_2 = (0.8, 0.8, 0.7, 0, 0.4, 0)^T, X_3 = (0.8, 0.8, 0.7, 0, 0, 0.4)^T$  and so on. (table 1)

In several lower solutions, all first component with same place look as if "tree-root", minimal component look as if "tree-tip", other look as if "tree-joint". Then lower Solution tree consist of all lower solution (fig. 2). The decimal in the fig.2 represents the component of lower solution, and integer in the parentheses represents inverse image, i.e. the place of above lower solution component.

The three lower solution trees are obtained from the above ten lower solutions. The supremum of each-class characteristic Fuzzy set is "tree-root", "tree-joint" and "tree-tip" separately. The branch of "lower Solution tree" is a lower solution, for example, the lower solution  $(0.8, 0.8, 0.7, 0, 0.4, 0)^T$  is represented by  $(0.8)_{(1,2)} \rightarrow (0.7)_{(1,3)} \rightarrow (0.4)_{(1,5)}$ .

Clearly, lower Solution tree is not equivalent class.

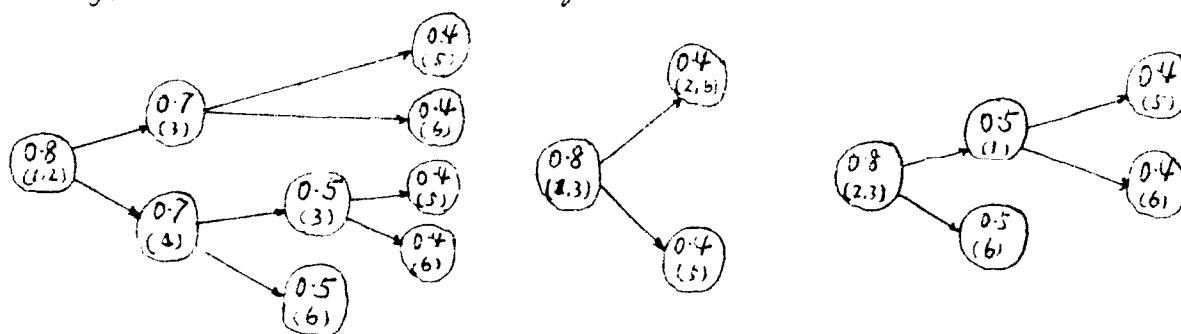


TABLE I

$R_{ij}$	Permutation	$i$	1	2	3	4	5	6	Lower solution at corresponding characteristic sum
$R_i$	$\frac{R_i}{1}$	$\frac{0.8}{1}$	$\frac{0.8}{1}$	$\frac{0.5}{6}$	$\frac{0.3}{9}$	$\frac{0.8}{1}$	$\frac{0.8}{3}$	$\frac{0.4}{7}$	$(0.8, 0.8, 0.5, 0.4, 0.3)$
Characteristic Sum									
1	$\lambda_1 = 0.8$		0.8	0.8	X	X	0.8	0.7	$(0.8, 0.8, X, X, X)$
2	$\lambda_1 + \lambda_2 = 0.8 + 0.7$					X			$(0.8, 0.8, X, 0.7, X, X)$
3	$\lambda_1 + \lambda_2 + \lambda_3 = 0.8 + 0.7 + 0.5$					0.5			$(0.8, 0.8, 0.5, 0.7, X, X)$
4	$\lambda_1 + \lambda_2 + \lambda_3 + \lambda_4 = 0.8 + 0.7 + 0.5 + 0.4$						0.4		$(0.8, 0.8, 0.5, 0.7, 0.4, 0)$
5	$0.8 + 0.7 + 0.5 + 0.4$		0.8	0.8	0	0	0	0.5	$(0.8, 0.8, 0.5, 0.7, 0.4, 0)$
6	$0.8 + 0.7 + 0.4$		0.8	0.8	0	0	0	0.7	$(0.8, 0.8, 0.7, 0, 0.4, 0)$
7	$0.8 + 0.7 + 0.4$		0.8	0.8	0	0	0	0.7	$(0.8, 0.8, 0.7, 0, 0.4, 0)$
8	$0.8 + 0.7 + 0.5$		0.8	0.8	0	0	0	0	$(0.8, 0.8, 0.7, 0, 0.5)$
9	$0.8 + 0.7 + 0.5 + 0.4$		0.8	0.8	0	0.8	0.8	0	$(0.8, 0.8, 0.5, 0.7, 0, 0.4)$
10	$0.8 + 0.4$		0.8	0.8	0	0	0	0.4	$(0.8, 0.8, 0.4, 0, 0.4, 0)$
11	$0.8 + 0.4$		0.8	0.8	0	0	0	0	$(0.8, 0.8, 0.4, 0, 0, 0)$
12	$0.8 + 0.5 + 0.4$		0	0	0.5	0	0.8	0	$(0.5, 0.8, 0.8, 0, 0.4, 0)$
13	$0.8 + 0.5 + 0.4$		0	0	0.5	0	0.8	0	$(0.5, 0.8, 0.8, 0, 0, 0.4)$
14	$0.8 + 0.5$		0	0	0	0.8	0.8	0	$(0, 0.8, 0.8, 0, 0, 0.5)$

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