

A STRUCTURAL REPRESENTATION OF THE TOTALLY FUZZY SETS.

by WU TAO

Many works have been done about the fuzzy sets, particularly about the fuzzy sets with a relation of indiscernability, called "the totally fuzzy sets". Applying the definition of the category of the totally fuzzy sets exposed by PONASSE (1), we will give a way of showing in order to represent the structure of a totally fuzzy set.

Thereafter, J is always designed by a finite sequence :

$$J = (r_0, r_1, \dots, r_n) \text{ with } 0 = r_0 < \dots < r_n = 1.$$

We recall the essential definition of JTF : the category of the J -totally fuzzy sets :

.The objects of JTF are the J -totally fuzzy sets, i-e the triplets $\underline{A} = (A, \alpha, \sigma)$ where A is a set, $\sigma: A \times A \rightarrow J$, $\alpha: A \rightarrow J$ are two functions verifying :

$$\sigma(a, a) = \alpha(a), \quad \sigma(a, b) = \sigma(b, a), \quad \sigma(a, b) \wedge \sigma(b, c) \leq \sigma(a, c)$$

. The morphisms from $\underline{A} = (A, \alpha, \sigma)$ to $\underline{B} = (B, \beta, \tau)$ are the binaries relations R from A to B , verifying :

- 1. $Ra \neq \emptyset$ where $Ra = \{b \in B \mid aRb\}$
- 2. aRb and $a'Rb' \implies \sigma(a, a') \leq \tau(b, b')$
- 3. aRb and $\alpha(a) \leq \tau(b, b') \implies aRb'$

. The composition of two morphisms

$$\underline{A} = (A, \alpha, \sigma) \xrightarrow{R} \underline{B} = (B, \beta, \tau) \xrightarrow{S} \underline{C} = (C, \nu, \rho)$$

is the morphism $SR: \underline{A} \rightarrow \underline{C}$ defined by :

$$aSRc \text{ iff there exists } b \in B, c' \in C \text{ such } aRb, bSc' \text{ and } \alpha(a) \leq \rho(c, c').$$

The identity $1_A: \underline{A} \rightarrow \underline{A}$ is $a1_A a' \iff \alpha(a) = \sigma(a, a')$.

. $R: \underline{A} \rightarrow \underline{B}$ is a monomorphism iff :

$$aRb \text{ and } a'Rb' \implies \sigma(a, a') = \alpha(a) \wedge \alpha(a') \wedge \tau(b, b').$$

G.MYCEK (3) has proved that the category JTF is a topos.

1. Lemma (4) : Given a TFS $\underline{A} = (A, \alpha, \sigma)$, if $a_0, a_1 \in A, a_0 \neq a_1$ and $\alpha(a_0) = \sigma(a_0, a_1)$ (i-e, $a_0 1_A a_1$) then we can "forget" a_0 . Precisely : let $\underline{A}' = (A', \alpha', \sigma')$ where $A' = A - \{a_0\}$ and $\sigma' = \sigma|_{A' \times A'}$ then \underline{A} is isomorphic to \underline{A}' in JTF.

Definition : Given $\underline{A} = (A, \alpha, \sigma)$, a frame of \underline{A} is an object $\underline{A}' = (A', \sigma', \sigma')$ of JTF such that \underline{A}' is isomorphic to \underline{A} and for any $a, a' \in A'$, $a \neq a' \implies \sigma'(a, a') < \alpha'(a) \wedge \alpha'(a')$.

2. Proposition : for any \underline{A} there exist a frame.

The demonstration is in (4). With the axiom of choice, the frame play the role of cardinals number of set.

3. Lemma : given $\underline{A} = (A, \alpha, \sigma)$ for $a_1, a_2, a_3 \in A$, let $r_1 = \sigma(a_2, a_3)$ $r_2 = \sigma(a_1, a_3)$ $r_3 = \sigma(a_1, a_2)$ and suppose that $r_1 \leq r_2 \leq r_3$, then $r_1 = r_2$.

4. Example : $J = (0, 0.2, 0.4, 0.6, 0.8, 1)$. $A = \{a, b, c, d, e, f, g, h, i, j, k\}$. To recognize the structure of \underline{A} it is enough to give a table of $\sigma(x, y)$, $x, y \in A$, a table of 11 x 11 square.

a	1										
b	0.8	1									
c	0.4	0.4	0.8								
d	0.4	0.4	0.6	0.6							
e	0.2	0.2	0.2	0.2	0.4						
f	0	0	0	0	0	0.6					
g	0	0	0	0	0	0.4	0.8				
h	0	0	0	0	0	0.2	0.2	1			
i	0	0	0	0	0	0	0	0	0.6		
j	0	0	0	0	0	0	0	0	0.6	0.6	
k	0	0	0	0	0	0	0	0	0.2	0.2	1
	a	b	c	d	e	f	g	h	i	j	k

by reason of symetry, it is sufficient to replenish the inferior triangle. For the first column we can choose the value freely under the condition $\sigma(a, x) \leq \alpha(a) \wedge \alpha(x)$, $x \in A$, but when we change the order of elements of A , it is convenient to choose a decreasing sequence. By the lemma 3, the second column is a reproduction of the first one. For the third column, $\sigma(c, d)$ may be chosen between 0.4 and 0.6, and for the rest we reproduce the second column. The sixth column, $\sigma(f, x)$ is freely chosen. We can certainly choose it decreasing. The ninth column, $\sigma(i, x)$ is freely chosen, and all the other columns reproduce the precedent ones.

So we see that in \underline{A} there are three groups of elements of A which are relatively closed to one another : $\{a, b, c, d, e\}$, $\{f, g, h\}$, $\{i, j, k\}$

and furthermore, among $\{a,b,c,d\}$ we see that a,b are close one another, **and** c,d are close.

5. Definition : given $\underline{A} = (A, \alpha, \sigma)$ and $r \in J$ $A_1 \subset A$, we said that A_1 is a indiscernible r -classe iff $A_1 \neq \emptyset$ and for all $a_1, a'_1 \in A_1$, $\sigma(a_1, a'_1) \geq r$ for all $a_1 \in A, a \in A - A_1 \implies \sigma(a_1, a) < r$.

6. Definition : let $r \in J$, we say that $\underline{A} = (A, \alpha, \sigma)$ is a r -direct sum of $\underline{A}_1 = (A_1, \alpha_1, \sigma_1)$ and $\underline{A}_2 = (A_2, \alpha_2, \sigma_2)$ and we note $\underline{A} = \underline{A}_1 \oplus_r \underline{A}_2$ if we have:

$$\bullet A = A_1 \sqcup A_2 \text{ (i.e } A_1 \subset A, A_2 \subset A, A_1 \cap A_2 = \emptyset, A_1 \cup A_2 = A, A_1 \neq \emptyset, A_2 \neq \emptyset \text{.)}$$

$$\bullet \sigma_i = \sigma|_{A_i \times A_i} \quad i = 1, 2 \quad \text{and:}$$

• For every $a_1 \in A_1, a_2 \in A_2, \sigma(a_1, a_2) = r$.

Remark : we can say that \underline{A}_1 and \underline{A}_2 are two r -disjoints sub-objets of \underline{A} and if $a_i, a'_i \in A_i$ ($i = 1, 2$), by lemma 3. we have

$$\sigma(a_i, a'_i) = \sigma_i(a_i, a'_i) \geq r.$$

In an other way, \underline{A}_i are two r -indiscernible classes.

7. Theorem : all J -totally fuzzy sets can be perfectly decomposed in r -direct sum, $r \in J$.

Demonstration : Let $\underline{A} = (A, \alpha, \sigma)$ and $L_r(\underline{A}) = \{a \in A \mid \alpha(a) \geq r\}$, $r \in J$. For each $r \in J$ such as $L_r(\underline{A}) \neq \emptyset$ the function $\sigma : A \times A \rightarrow J$ inducing an equivalent relation $L_r(\sigma)$ on $L(\underline{A})$:

$$a L_r(\sigma) a' \iff \sigma(a, a') > r.$$

Remarking that $J = \{0 = r_0 < r_1 < r < \dots < r_n = 1\}$. If $\underline{A} \neq \emptyset$, we can certainly suppose that $L_{r_1}(\underline{A}) = A$. So the equivalent relation $L_{r_1}(\sigma)$ determines a partition on $A = L_{r_1}(\underline{A})$:

$A = \bigsqcup_{i \in I} A_i^{r_1}$, if $a_i \in A_i^{r_1}, a_j \in A_j^{r_1}$ with $i \neq j$, We have $\sigma(a_i, a_j) < r_1$ so $\sigma(a_i, a_j) = 0$. So we obtain a 0-direct sum : $\underline{A} = \bigoplus_{i \in I} \underline{A}_i^{r_1}$. More clearly, we apply a pair of parenthesis of level 0 : ${}^0(\ ,)^0 : i \in I$.

$$\underline{A} = {}^0\left(\bigoplus_{i \in I} \underline{A}_i^{r_1}\right)^0$$

where each $\underline{A}_i^{r_1}$ is a r_1 -indiscernible classe, so we say that \underline{A} is decomposed in 0-sum of his r_1 -indiscernible class. We **return** take the $\underline{A}_i^{r_1}$ for \underline{A} : for each $\underline{A}_i^{r_1}$ such as $L_{r_1}(\underline{A}_i^{r_1}) \neq \emptyset$, the equivalent relation $L_{r_2}(\sigma)$ determin a partition on $L_{r_1}(\underline{A}_i^{r_1})$:

$L_{r_1}(\underline{A}_i^{r_1}) = \bigsqcup_{j \in J_i} A_{i,j}^{r_2}$ We have for all $a_j, a'_j \in A_{i,j}^{r_2}$ $\sigma(a_j, a'_j) > r_1$ and for all

$$a_j \in A_{i,j}^{r_2}, \quad a_{j'} \in A_{i,j'}^{r_2}, \quad j \neq j' \quad (a_j, a_{j'}) = r_1.$$

By the same procedure the r_1 -indiscernible class $A_i^{r_1}$ is a r_1 -direct sum of its r_2 -indiscernible classes :

$$A_i^{r_1} = \bigoplus_{j \in J_i} A_{i,j}^{r_2} = \bigoplus_{j \in J_i} (A_{i,j}^{r_2})^1$$

and so we have :

$$A = \bigoplus_{i \in I} \left(\bigoplus_{j \in J_i} (A_{i,j}^{r_2})^1 \right)^1$$

Applying at most n times this procedure, we obtain a complete decomposition of A in all its indiscernible classes. For example, we take A as in example of 4. We have the following decomposition :

$$A = (((\underline{a} \oplus \underline{b}) \oplus (\underline{c} \oplus \underline{d})) \oplus \underline{e}) \oplus ((\underline{f} \oplus \underline{g}) \oplus \underline{h}) \oplus ((\underline{i} \oplus \underline{j}) \oplus \underline{k});$$

0.8 0.4 0.6 0.2 0 0.4 0.2 0 0.6 0.2

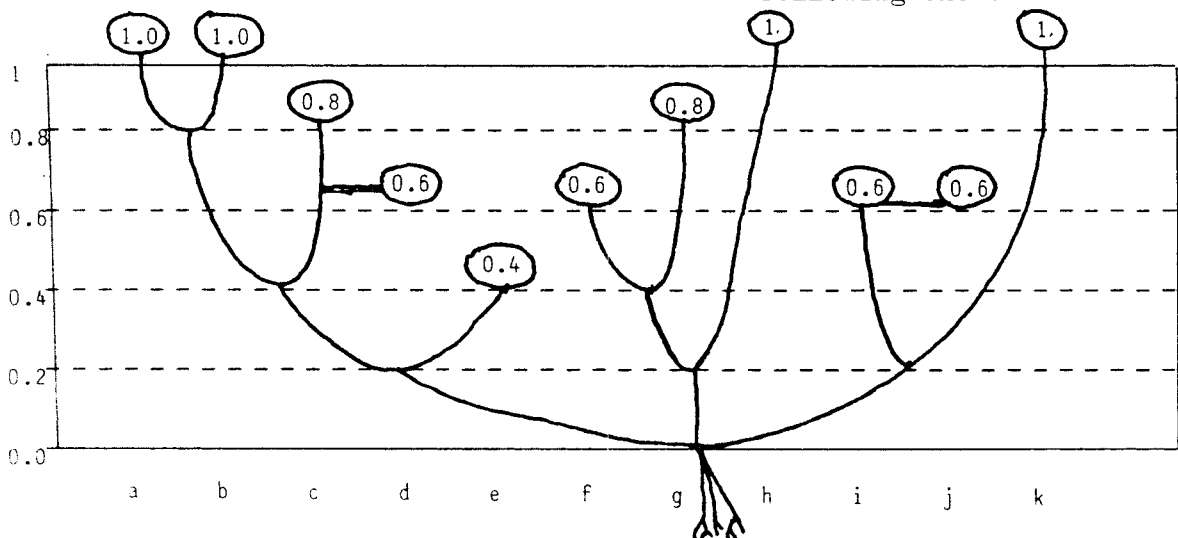
Moreover, it is possible to input this structural information of A in a computer in the form :

$$= (((((1)0.8(1))0.4((0.8)0.6(0.6))0.2(0.4))\text{---} \\ \text{---}(((0.6)0.4(0.8))0.2(1))\dots(((0.6)0.6(0.6))0.2(1)).$$

8. The tree of indiscernible decomposition of a A :

In the demonstration of the theorem 7 we saw that if $a, a' \in A$ $\sigma(a, a') = r_i$ there exist a unique pair of " r_i -level" parenthesis enclosing a and a' : $i(\dots a, \dots, a', \dots)^i$. so we can draw a tree where the elements of A are considered as leaves with the value $\alpha(x)$. Given two leaves a and a' , they are connected by one way by branches and $\sigma(a, a')$ is the value attributed to the lowest branch.

For example, take A as in the example of 4 : its tree is the following one :



We remark that :

.we can immediately recognize the structure of \underline{A} : $\alpha(x)$ and $\sigma(x,y)$.

.the elements we can forget are d and j (or d and i) because their are linked by horizontal branches.

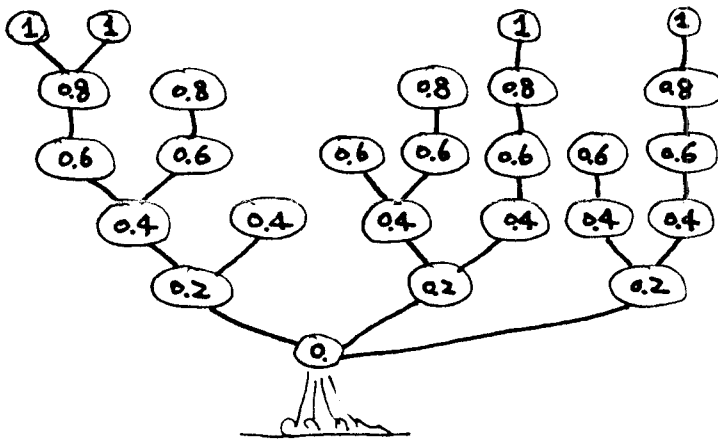
.the decomposition scheme clearly appears in the tree.

.we recognize immediately the r -indiscernible classes of A . For all $r \in J$.

.by the lemma 1, we can abandon d and j which are not necessary but we can also fullfill to each level a new leave on the branches. The lemma 1 permits this inverse proceeding : obtaining a full tree of \underline{A} ;

.According to the meaning of the tree we can explain many elementary

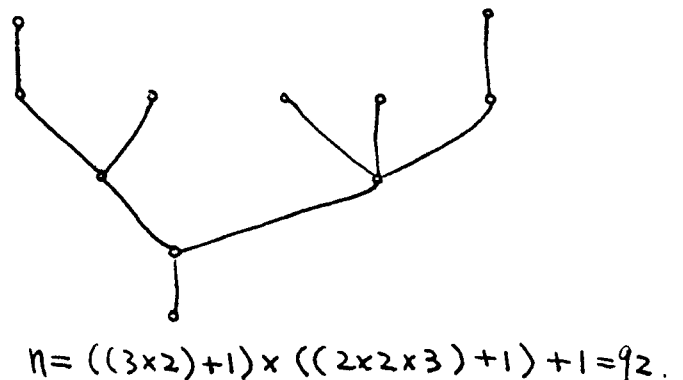
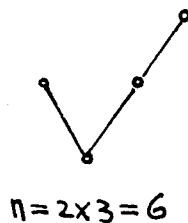
notions as morphism, sub-object of a total fuzzy set, equalizer, etc... For example, given \underline{A} , let $\mathcal{C}_A = \{r\text{-indiscernible classe of } \underline{A} \mid r \in J\}$, we have a morphism $R : \underline{A} \rightarrow \underline{B}$ in bijection with a function $\tilde{R} : \mathcal{C}_A \rightarrow \mathcal{C}_B$ such that if $x \in \mathcal{C}$ is a r -classe, so $y = \tilde{R}(x) \in \mathcal{C}_B$ is also a r -classe. If $y' = \tilde{R}(x')$ then we have $\sigma(x,x') \leq \tau(y,y')$ (σ is the prolongation of σ from A to $\mathcal{C}_A : A \subset \mathcal{C}_A$ since each $a \in A$ is a $\alpha(a)$ -classe



the full tree of \underline{A} .

We also can calculate the number n of sub-objects of \underline{A} : the tree is a powerful idea to study the structure of any abstract category (4);

. According to the meaning of the tree of \underline{A} we can calculate the number n of all sub-objects of \underline{A} :



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