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Let  $G$  be a given set,  $W$  be a mathematical system. If  $g \in W \uparrow G^2$ ,  $g$  is called  $W$ -binary relation of  $G$ .

Pansystems Reduce Principle<sup>[1]</sup>: Let  $g \subset G^2 \times W$ ,  $D \subset W$ , then  $g \circ D \subset G^2$ ,  $G = \bigcup G_i (d \delta (f))$ , where  $\delta$  is a pansystems operator,  $f \in \{(g \circ D \vee I)^a, (g \circ D)^b\}$   $a, b = (n), t, n = 0, 1, \dots$ .

The principle provides more than one hundred theorems to dispose of generalized fuzzy network. The generalized fuzzy network is a network with panweight. Pansystems network analysis is the generalized fuzzy network studied by viewpoint and method of pansystems methodology. It mainly investigates series-parallel analysis, clustering-discoupling, shengke automaton network. In this paper, the so-called flow is discussed by viewpoint of pansystems methodology.

At first, we introduce the concepts of the quasi-opti semi-field and  $\theta$ -composition.

Definition 1<sup>[1]</sup> ( $\theta$ -composition) Let  $\theta_1 \in W \uparrow P(W)$ ,  $\theta_2 \in W \uparrow W^2$ .  
 $\theta = (\theta_1, \theta_2): (W \uparrow G^2)^2 \rightarrow W \uparrow G^2$ ,  $(g_1 \theta g_2)(x, y) = \theta_1 \{ g_1(x, t) \theta_2 g_2(t, y) \mid t \in G \}$ .

Definition 2 If  $\theta_1, \theta_2$  satisfy the following conditions,  $(W, \theta_1, \theta_2)$  is called a quasi-opti-semi-field.

- 1 association and distribution law .
- 2  $\theta_1(a, b) = a$  or  $\theta_1(a, b) = b$  .
- 3 there exists  $0 \in W$ ,  $\theta_2(a, 0) = a$  for any  $a \in W$  .

If  $f, g \in W \uparrow G^2$ ,  $\theta_1(f(x, y), g(x, y)) = g(x, y)$ , denote  $f \leq g$ .  
 Let  $F = \{f \mid g_1 \leq f \leq g_2, f_0 \bar{\theta} f(x, x) = f \bar{\theta} f_0(x, x), f \in W \uparrow G^2\}$  where  $f_0 \in \{0\} \uparrow G^2$ ,  $g_i \in W \uparrow G^2$   $i=1, 2$ .  $x \in W$ ,  $f \bar{\theta} g(x, y) = \theta_2 \{f(x, t) \theta_1 g(t, y) \mid t \in G\}$   $f \in F$  is called a panweight flow of network  $(G, g_1, g_2)$ .

Theorem 1  $f \in W \uparrow G^2$  is a panweight flow if and only if  $g_1 \leq f \leq g_2$  and  $f \bar{\theta} f_0 = f_0 \bar{\theta} f^{-1}$ .

Theorem 2 If  $F \neq \phi$ , then  $g_1 \leq f_0 \bar{\theta} g_2^m$ ,  $g_1 \leq g_2^m \bar{\theta} f_0$  for any integer  $m$ , where  $g^m = g \bar{\theta} g^{m-1}$ . and  $\theta_1, \theta_2$  satisfy the condition: For  $a, b, c \in W$  with  $\theta_1(a_i, b_i) = b_i$ ,  $i=1, 2$ , then  $\theta_1(\theta_2(a_1, a_2), \theta_2(b_1, b_2)) = \theta_2(b_1, b_2)$ .

For network  $(G, g_1, g_2)$ , let  $G = \{s\} \cup G_1$ ,  $g_1(x, s) = 0$ , for any  $x \in G_1$ , and define  $V \in W \uparrow F$ ,  $V(f) = f \circ \bar{0} \circ f_0(s, s)$ ,

$$E_f = \{f \mid V(f_i) \neq V(f) \text{ and } \theta_1(V(f), V(f_i)) = V(f_i), f \in F\}$$

$$E_f^* = \{f \mid 1 \neq V(f) \text{ and } \theta_1(1, V(f)) \neq 1 \in V(F)\}, \text{ where } f_i \in F.$$

$$\text{Theorem 3 } |E_f| \geq |E_f^*|, V(E_f) = E_f^*.$$

From the equivalence corresponding principle<sup>[1]</sup>,  $V \circ V^{-1} \in E[F]$ .

Theorem 4 If  $f_i \in E_i \subset E_i(d(V \circ V^{-1}))$  such that  $|E_i| = 0$ ,

then: (i)  $\bigcup_{f \in E_i} f = \emptyset$ , (ii)  $E_i \subset E_f$  for any  $f \in F - E_i$ .

Let  $V^* \in W \uparrow Q$  be induced by  $V$ ,  $V(f_i) = V^*(E_i)$ , where  $Q = F/V \circ V^{-1}$ ,  $f_i \in E_i$ . Now define a binary relation  $H \in P(Q^2)$ ,  $(E_i, E_j) \in H$  if and only if  $\theta_1(V^*(E_i), V^*(E_j)) = V^*(E_j)$ .

Theorem 5 (i)  $H \in L_S(Q)$ , (ii)  $(E_i, E_j) \in H$  if and only if  $E_i \subset E_j$ , where  $f_i \in E_i$ ,  $f_j \in E_j$ .

Pansystems reduce principle plays an important role in treating with series-parallel, clustering-discoupling. We apply it to pansystems network analysis.

Let  $(N, +, V)$  be a integer quasi-opti-semi-field, where  $a \vee b = \max\{a, b\}$ . Network  $(G, g_1, g_2)$ , where  $g_1 \in \{0\} \uparrow G^2$ ,  $s, t$  be source and sink respectively.  $f$  be a flow. Define

$$N_f = \{n \mid \text{there exists } (x, y) \in g_2 \circ n, \text{ such that } (x, y) \circ g_2 > (x, y) \circ f\}.$$

Theorem 6 If  $G_i \subset G$  ( $d \xi_i(f_i^*)$ ),  $f_i = g_2 \circ N$ , then  $f$  is a maximal flow if and only if  $(s, t) \bar{\in} G_i$  for any  $i$ .

#### References

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